

DYNAMICS GROUPS OF ASYNCHRONOUS CELLULAR AUTOMATA

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ABSTRACT. A *sequential dynamical system*, or *SDS*, consists of an undirected graph Y , a vertex-indexed list of update rules \mathfrak{F}_Y , and a word ω over the vertex set that contains each vertex at least once and describes the order in which these update rules are to be applied. When an SDS is ω -*independent*, i.e., when the set of periodic states does not depend on the choice of ω , then the update rules permute the periodic states and these permutations generate the *dynamics group* of the SDS. In this article we investigate dynamics groups in the special case where Y is a circular graph with n vertices, there are only two vertex states, and all of the update rules are identical. The 256 possibilities are known as *Wolfram rules* and the resulting sequential dynamical systems are called *asynchronous cellular automata*, or ACAs. We previously showed that for all $n > 3$ exactly 104 of the 256 Wolfram rules give rise to ω -independent ACAs. Here we classify the periodic states of these ω -independent systems and describe their dynamics groups, which are quotients of Coxeter groups.

A cellular automaton, or CA, is a classical discrete dynamical system defined over a regular grid of cells. Every cell takes on one of a finite number of states and has an update rule that only depends on its state and the states of its neighbors. At every discrete time step, the update rules are simultaneously applied. More recent work has investigated *sequential dynamical systems* defined over arbitrary finite graphs where the update rules are applied asynchronously. (See [3] for a detailed bibliography.) The special case where Y is a circular graph, there are only two vertex states, and the update rules are all the same is the class of elementary *asynchronous cellular automata* or ACAs. There are 256 possible update rules. In an earlier article ([2]) we proved that exactly 104 of these give rise to an ACA whose periodic states are independent of the update order, a property independent of the size of the underlying graph. In this article we describe the periodic states and the dynamics group for each ω -independent ACA identified in [2].

The article is structured as follows. Sections 1, 2 and 3 contain background definitions, notations, and results. Section 4 studies ACAs whose dynamics groups are trivial, Section 5 focuses on Wolfram rules that are invertible, and Section 6 investigates the remaining cases. The results are summarized in Tables 1, 3, and 5.

1. SEQUENTIAL DYNAMICAL SYSTEMS

A *sequential dynamical system*, or *SDS*, is a discrete dynamical system with three components: an undirected graph Y , a list of update rules \mathfrak{F}_Y , and an update order ω .

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Definition 1.1 (Graph conventions). Let Y be a simple undirected graph with n vertices labeled from 1 to n , and recall that the *neighbors* of a vertex are those vertices connected to it by an edge. If \mathbb{F} is a finite field and every vertex is assigned a value from \mathbb{F} , then a global state of the system is described by an n -tuple \mathbf{y} whose i^{th} coordinate indicates the current state of the vertex i . The set of all possible states is the vector space \mathbb{F}^n .

Definition 1.2 (Local functions). A function $F: \mathbb{F}^n \rightarrow \mathbb{F}^n$ is called *Y -local at i* if for each $\mathbf{y} \in \mathbb{F}^n$ (1) $F(\mathbf{y})$ only alters the i^{th} coordinate of \mathbf{y} and (2) the new value of the i^{th} coordinate only depends on the coordinates of \mathbf{y} corresponding to i and its neighbors in Y . Other names for such a function are a *local function* or an *update rule*. We use \mathfrak{F}_Y to denote a list with one local function for each vertex of Y . More precisely, $\mathfrak{F}_Y = (F_1, F_2, \dots, F_n)$ where F_i is a function that is Y -local at i .

Definition 1.3 (Restricted local functions). If i is a vertex with k neighbors in Y , then corresponding to each function F that is Y -local at i , we define a function $f: \mathbb{F}^{k+1} \rightarrow \mathbb{F}$ where the domain is restricted to the coordinates corresponding to i and its neighbors, and the output is the new value F would assign to the i^{th} coordinate under these conditions. The functions F and f contain the same information packaged differently and each determines the other. Both have their uses. Functions such as F can be readily composed, but functions such as f are easier to explicitly describe.

Definition 1.4 (Update orders). An *update order* ω is a finite sequence of numbers chosen from the set $\{1, \dots, n\}$. If every number $1 \leq i \leq n$ occurs at least once, we say it is *fair* and if every number occurs exactly once, then it is *simple*. We use the notation $\omega = (\omega_1, \omega_2, \dots, \omega_m)$ with $m = |\omega|$. Of course, $m \geq n$ when ω is fair and $m = n$ when ω is simple. Let W_Y denote the collection of all update orders and let S_Y denote the subset of simple update orders. The subscript Y indicates that we are thinking of the numbers in these sequences as vertices in the graph Y .

Definition 1.5 (Sequential dynamical systems). A *sequential dynamical system*, or *SDS*, is a triple $(Y, \mathfrak{F}_Y, \omega)$ consisting of an undirected graph Y , a list of local functions \mathfrak{F}_Y , and a fair update order $\omega \in W_Y$. If ω is the sequence $(\omega_1, \omega_2, \dots, \omega_m)$, then we construct the *SDS map* $[\mathfrak{F}_Y, \omega]: \mathbb{F}^n \rightarrow \mathbb{F}^n$ as the composition $[\mathfrak{F}_Y, \omega] := F_{\omega_m} \circ \dots \circ F_{\omega_1}$.

The main goal is to understand the dynamics of the SDS map, i.e., its behavior under iteration.

Definition 1.6 (ω -independence). Let $\text{Per}[\mathfrak{F}_Y, \omega]$ denote the set of states periodic under iterations of $[\mathfrak{F}_Y, \omega]$. A list of Y -local functions \mathfrak{F}_Y is called *ω -independent* if $\text{Per}[\mathfrak{F}_Y, \omega] = \text{Per}[\mathfrak{F}_Y, \omega']$ for all fair update orders $\omega, \omega' \in W_Y$ and *π -independent* if $\text{Per}[\mathfrak{F}_Y, \pi] = \text{Per}[\mathfrak{F}_Y, \pi']$ for all simple update orders $\pi, \pi' \in S_Y$. When \mathfrak{F}_Y is ω -independent, we write $\text{Per}(\mathfrak{F}_Y)$ instead of $\text{Per}[\mathfrak{F}_Y, \omega]$.

In the case of ω -independence we may – by abuse of notation – let the list of Y -local functions \mathfrak{F}_Y stand for the entire SDS. When \mathfrak{F}_Y is ω -independent, its local functions permute the elements of $\text{Per}(\mathfrak{F}_Y)$.

Proposition 1.7 (Permuting periodic states). *If \mathfrak{F}_Y is ω -independent and $P = \text{Per}(\mathfrak{F}_Y)$, then for each i , $F_i(P) = P$. In particular, the restriction of F_i to P is a permutation.*

Proof. Let $\omega = (\omega_1, \omega_2, \dots, \omega_m)$ be a fair update order with $\omega_1 = i$ and let σ be the modified update order with ω_1 moved from the first to last: $\sigma = (\omega_2, \omega_3, \dots, \omega_m, \omega_1)$. Since $F_i \circ [\mathfrak{F}_Y, \omega]^k = [\mathfrak{F}_Y, \sigma]^k \circ F_i$ for all k , and by hypothesis $[\mathfrak{F}_Y, \omega]^k(\mathbb{F}^n) = [\mathfrak{F}_Y, \sigma]^k(\mathbb{F}^n) = P$ for all sufficiently large k , we find that $F_i(P) \subset P$. More explicitly, for large enough k ,

$$F_i(P) = F_i \circ [\mathfrak{F}_Y, \omega]^k(\mathbb{F}^n) = [\mathfrak{F}_Y, \sigma]^k \circ F_i(\mathbb{F}^n) \subset P.$$

Moreover, since $[\mathfrak{F}_Y, \omega](P) = P$, the restriction of F_i to P is injective, and $F_i(P) = P$. \square

When \mathfrak{F}_Y is ω -independent, we write F_i^* and $[\mathfrak{F}_Y, \omega]^*$ to denote the restrictions of these maps to $\text{Per}(\mathfrak{F}_Y)$. By Proposition 1.7 all such maps are permutations. Note that ω -independence focuses on the periodic states as a set rather than how these states are permuted. In particular, when an ω -independent \mathfrak{F}_Y is paired with two different update orders ω and σ , the permutations $[\mathfrak{F}_Y, \omega]^*$ and $[\mathfrak{F}_Y, \sigma]^*$ are often distinct. These various permutations can be used to construct a group encoding all of the possible dynamics [3].

Definition 1.8 (Dynamics group). Let \mathfrak{F}_Y be ω -independent. For any collection of update orders $U \subseteq W_Y$, the *dynamics group of \mathfrak{F}_Y with respect to U* is

$$\text{DG}(\mathfrak{F}_Y, U) = \langle [\mathfrak{F}_Y, \omega]^* \mid \omega \in U \rangle.$$

It should be clear that when U and V are sets of update orders and U is contained in the closure of V under concatenation, then $\text{DG}(\mathfrak{F}_Y, U) \subset \text{DG}(\mathfrak{F}_Y, V)$. The *dynamics group of \mathfrak{F}_Y* , $\text{DG}(\mathfrak{F}_Y) = \text{DG}(\mathfrak{F}_Y, W_Y)$, and the *restricted dynamics group of \mathfrak{F}_Y* , $\text{RDG}(\mathfrak{F}_Y) = \text{DG}(\mathfrak{F}_Y, S_Y)$, are special cases of particular interest. Note that $\text{DG}(\mathfrak{F}_Y)$ contains and is generated by the bijections F_i^* .

2. ASYNCHRONOUS CELLULAR AUTOMATA

The SDSs we focus on are defined over circular graphs, they have only two possible vertex states, and all the local functions are identically defined. These *asynchronous cellular automata*, or ACAs, are an asynchronous version of the classical finite elementary cellular automata. Even in such a restrictive situation there are many interesting dynamical behaviors.

Definition 2.1 (Circular graphs and vertex states). Let $Y = \text{Circ}_n$ denote the *circular graph* with vertex set $\{1, \dots, n\}$ (viewed as residue classes mod n) and edges connecting i and $i + 1 \pmod n$. To avoid trivialities, we always assume $n > 3$. Each vertex has two possible states that we identify with $\mathbb{F}_2 = \{0, 1\}$, the field of size 2.

Definition 2.2 (Wolfram rules). Let $F_i: \mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ be a function Circ_n -local at i and let $f_i: \mathbb{F}_2^3 \rightarrow \mathbb{F}_2$ be its restricted form. Because the neighbors of i are $i - 1$ and $i + 1$, it is conventional to list these coordinates in ascending order in the domain of f_i , keeping in mind that all subscripts are viewed mod n . The function F_i , henceforth referred to as a *Wolfram rule*, updates the value of y_i based on the value of the triple (y_{i-1}, y_i, y_{i+1}) and

it is completely determined by how the i^{th} coordinate is updated in these eight possible situations. In other words, F_i is completely described by the following table.

$y_{i-1}y_iy_{i+1}$	111	110	101	100	011	010	001	000
$f_i(y_{i-1}, y_i, y_{i+1})$	a_7	a_6	a_5	a_4	a_3	a_2	a_1	a_0

More concisely, the $2^8 = 256$ possible Wolfram rules can be indexed by an 8-digit binary number $a_7a_6a_5a_4a_3a_2a_1a_0$, or by its decimal equivalent $k = \sum a_i2^i$. There is thus a *Wolfram rule* k for each integer $0 \leq k \leq 255$.

Definition 2.3 (Asynchronous cellular automata). We write $\mathbf{Wolf}_i^{(k)}$ to denote the update rule $F_i: \mathbb{F}^n \rightarrow \mathbb{F}^n$ corresponding to k and $\mathfrak{Wolf}_n^{(k)}$ for the list $(\mathbf{Wolf}_1^{(k)}, \mathbf{Wolf}_2^{(k)}, \dots, \mathbf{Wolf}_n^{(k)})$ of update rules. For each fair update order ω the SDS $(\text{Circ}_n, \mathfrak{Wolf}_n^{(k)}, \omega)$ is called an *asynchronous cellular automaton*, or *ACA*. If $\mathfrak{Wolf}_n^{(k)}$ is ω -independent for all $n > 3$, we say Wolfram rule k is ω -independent.

When $\mathfrak{Wolf}_n^{(k)}$ is ω -independent, let $P_{n,k} = \text{Per}(\mathfrak{Wolf}_n^{(k)})$ denote its periodic states and let $G_{n,k} = \text{DG}(\mathfrak{Wolf}_n^{(k)})$ denote its dynamics group. We usually suppress the dependence on n and write simply P_k and G_k . In this notation, our goal is to describe the set P_k and the group G_k for each ω -independent Wolfram rule. In [2] we proved the following result.

Theorem 2.4. *Exactly 104 Wolfram rules are ω -independent. More precisely, $\mathfrak{Wolf}_n^{(k)}$ is ω -independent for all $n > 3$ iff $k \in \{0, 1, 4, 5, 8, 9, 12, 13, 28, 29, 32, 40, 51, 54, 57, 60, 64, 65, 68, 69, 70, 71, 72, 73, 76, 77, 78, 79, 92, 93, 94, 95, 96, 99, 102, 105, 108, 109, 110, 111, 124, 125, 126, 127, 128, 129, 132, 133, 136, 137, 140, 141, 147, 150, 152, 153, 156, 157, 160, 164, 168, 172, 184, 188, 192, 193, 194, 195, 196, 197, 198, 199, 200, 201, 202, 204, 205, 206, 207, 216, 218, 220, 221, 222, 223, 224, 226, 228, 230, 232, 234, 235, 236, 237, 238, 239, 248, 249, 250, 251, 252, 253, 254, 255\}$.*

In [2] we also defined the inversion, reflection, and inversion-reflection of an ACA. Loosely speaking, *inversion* systematically swaps the roles of 0 and 1, *reflection* systematically switches left and right, and *inversion-reflection* does both at once. The inversion, reflection, or inversion-reflection of an ω -independent Wolfram rule is still ω -independent, it has a corresponding set of periodic states and an isomorphic dynamics group. This should not be surprising since all we have done is relabel the underlying states on which the local functions act. Rules related in this manner are *dynamically equivalent*. When the 256 Wolfram rules are partitioned into classes of rules related by reflection, inversion or both, there are 88 equivalence classes. The 104 rules listed in Theorem 2.4 belong to 41 such classes and thus we only need to describe P_k and G_k for 41 representative values of k .

Corollary 2.5 (41 representative rules). *Every ω -independent Wolfram rule is dynamically equivalent to $\mathfrak{Wolf}_n^{(k)}$ for some $k \in \{0, 1, 4, 5, 8, 9, 12, 13, 28, 29, 32, 40, 51, 54, 57, 60, 72, 73, 76, 77, 105, 128, 129, 132, 133, 136, 137, 140, 141, 150, 152, 156, 160, 164, 168, 172, 184, 200, 201, 204, 232\}$.*

Finally, because the binary notation is cumbersome and the decimal notation is opaque, we introduce (as in [2]) a concise symbolic tag for each Wolfram rule.

Definition 2.6 (Tags). The four functions from \mathbb{F}_2 to \mathbb{F}_2 can be described by their behavior: the value never changes, the value always changes, both elements go to 0, or both elements go to 1. We refer to these functions by the evocative symbols $-$, \mathbf{x} , 0 , and 1 , respectively. The restricted local form of Wolfram rule k is completely determined by the four functions from \mathbb{F}_2 to \mathbb{F}_2 that result when the values of y_{i-1} and y_{i+1} are held constant. Let t_0 , t_1 , t_2 and t_3 be the symbols for these functions in the four cases $y_{i-1} = y_{i+1} = 0$, $y_{i-1} = 0$ and $y_{i+1} = 1$, $y_{i-1} = 1$ and $y_{i+1} = 0$, and $y_{i-1} = y_{i+1} = 1$, respectively. The *tag* of k is the string $t_3t_2t_1t_0$. Note that t_0 depends on the values of a_0 and a_2 , t_1 depends on a_1 and a_3 , t_2 depends on a_4 and a_6 , and t_3 depends on a_5 and a_7 . The numbering and the order of the t_i 's has been chosen to match the traditional binary representation as closely as possible, easing the transition between the two. As an illustration, the reader can verify that Wolfram rule 29 has binary notation 00011101 and tag $0\mathbf{x}-1$.

Patterns among the Wolfram rules are easier to discern when using tags.

Remark 2.7 (Tags and dynamic equivalence). On the level of tags, reflections switch the order of t_1 and t_2 . For example, the reflection of rule $0-1\mathbf{x}$ is $01-\mathbf{x}$. To describe the effect that inversion has on tags, we define a map $\iota: \{1, 0, -, \mathbf{x}\} \rightarrow \{1, 0, -, \mathbf{x}\}$ that fixes $-$ and \mathbf{x} while switching 0 and 1 . When Wolfram rule k has tag $t_3t_2t_1t_0$, its inversion has tag $\iota(t_0)\iota(t_1)\iota(t_2)\iota(t_3)$. For example, the inversion of rule $0-1\mathbf{x}$ is $\mathbf{x}0-1$. See [2] for a more detailed explanation.

Remark 2.8 (Tags and other SDSs). If $(Y, \mathfrak{F}_Y, \omega)$ is an SDS with only two vertex states, then each update rule F_i can be described by a set of symbols similar to the tag used to describe Wolfram rules. More specifically, if vertex i has exactly k neighbors, then the restricted local form of F_i is a function $f_i: \mathbb{F}_2^{k+1} \rightarrow \mathbb{F}_2$ and this function is determined by the 2^k functions from \mathbb{F}_2 to \mathbb{F}_2 that result when the values of the neighbors of i are held constant. In particular, the behavior of F_i is determined by the corresponding 2^k symbols, selected from $\{1, 0, -, \mathbf{x}\}$.

3. GROUPS

In preparation for our investigation of dynamics groups we recall a few basic facts about group actions and Coxeter groups.

Definition 3.1 (Faithful actions). Let G be a group acting on a set X . The action is called *faithful* if nontrivial elements act nontrivially. When this is the case, we can view G as a subgroup of SYM_X . Note that groups generated by permutations act faithfully on their underlying sets almost by definition. In particular, for any ω -independent SDS \mathfrak{F}_Y , the action of its dynamics group $\text{DG}(\mathfrak{F}_Y)$ on its periodic states $\text{Per}(\mathfrak{F}_Y)$ is faithful.

One of the key features of a group acting on a set is its orbit structure.

Definition 3.2 (Orbits). Let G be a group acting on a finite set X . The *orbit* of a point $x \in X$ is the subset of points to which it can be sent by an element of G . Thus the orbit of x is $Gx = \{g \cdot x \mid g \in G\}$. Because two orbits are either identical or disjoint, the collection

of all orbits $\{Gx \mid x \in X\}$ partitions X into equivalence classes $X = X_1 \sqcup X_2 \sqcup \dots \sqcup X_\ell$. When there is only one orbit we say the action is *transitive*.

As we investigate the faithful action of $G_{n,k}$ on $P_{n,k}$, we use $o_{n,k}$ to denote the number of orbits under this action. And, as with G_k and P_k , we usually suppress the dependence on n and write o_k instead. The orbit structure is of interest because distinct orbits provide information about the structure of the group.

Definition 3.3 (Subdirect products). If G acts faithfully on a finite set X with orbits X_1, X_2, \dots, X_ℓ then G can be viewed as a subgroup of $\text{SYM}_{X_1} \times \text{SYM}_{X_2} \times \dots \times \text{SYM}_{X_\ell}$. To see this note that when an element of G is written in disjoint cycle notation, each cycle must permute elements within a single orbit. Thus every $g \in G$ can be viewed as a ℓ -tuple $g = (g_1, g_2, \dots, g_\ell)$ where $g_i \in \text{SYM}(X_i)$. Even better, we can replace each SYM_{X_i} with G_i , the image of G under the projection to the i^{th} factor. As a result G embeds in a direct product $G_1 \times G_2 \times \dots \times G_\ell$, where the projection to each factor is onto. We call G a *subdirect product* of G_1, G_2, \dots , and G_ℓ , and the G_i s are called the *orbit groups* of G .

For an illustration of these concepts, consider a group generated by a single permutation.

Example 3.4 (Cyclic groups). The permutation $(1, 2)(3, 4, 5)(6, 7, 8, 9)$ generates a cyclic group G of order 12. It naively belongs to SYM_9 , a group of size 362880, but based on its orbit structure it lives inside the much smaller group $\text{SYM}_2 \times \text{SYM}_3 \times \text{SYM}_4$ of size 288. And if we cut down each factor to the image under projection, then G embeds in $\mathbb{Z}_2 \times \mathbb{Z}_3 \times \mathbb{Z}_4$, a noncyclic group of order 24 that contains G as an index 2 subgroup.

One caution is that G can be a subdirect product of groups without splitting as a direct product. Consider the cyclic group G generated by $(1, 2, 3)(4, 5, 6)$. The procedure described above embeds G as a subdirect product of \mathbb{Z}_3 and \mathbb{Z}_3 , but G itself is simple. The second set of results we need to recall are about Coxeter groups. A Coxeter group is a group generated by involutions with a presentation of a particularly simple form.

Definition 3.5 (Coxeter groups). Let W be a group generated by a finite subset $S = \{s_1, \dots, s_n\}$ and let $m_{ij} \in \mathbb{N} \cup \{\infty\}$ denote the order of the product $s_i s_j$. The pair (W, S) is called a *Coxeter system* if the elements of S are involutions (i.e., $m_{ii} = 1$) and the presentation $\langle s_1, \dots, s_n \mid (s_i s_j)^{m_{ij}} = 1 \rangle$ is a presentation of W . Note that when $m_{ij} = \infty$, no relation is included for this pair of indices and that trivially $m_{ij} = m_{ji} \geq 2$ for all $i \neq j$.

The letters W and S are those traditionally used for a *Coxeter group* and its *Coxeter generators*. Coxeter groups have an incredibly rich theory and close connections with many areas of mathematics. By comparison, the results we need are fairly modest. Before listing these results, we first establish the relevance of Coxeter groups to our investigation of dynamics groups.

Proposition 3.6 (Coxeter quotients). *Every group generated by a finite set of involutions can be viewed as a quotient of a Coxeter group in a natural way.*

Proof. Let G be a group and let $S = \{s_1, \dots, s_n\}$ be a subset of involutions that generate G . If we define m_{ij} as the order of $s_i s_j$ in G and we define W as the group with presentation

$\langle s_1, \dots, s_n \mid (s_i s_j)^{m_{ij}} = 1 \rangle$ then there is a natural surjective homomorphism from W to G sending s_i to s_i . \square

Theorem 3.7 (Dynamics groups as Coxeter quotients). *If \mathfrak{F}_Y is an ω -independent SDS with only two possible vertex states, then F_i^* , the restriction of a local function to the periodic states $\text{Per}(\mathfrak{F}_Y)$, is either trivial or an involution. As a consequence, the dynamics group $\text{DG}(\mathfrak{F}_Y)$ is either trivial or a quotient of a Coxeter group.*

Proof. The key observation is that because F_i^* can only change the i^{th} coordinate, the size of the cycles in its cycle structure are bounded by the number of possible vertex states. For the second assertion, note that when the dynamics group $\text{DG}(\mathfrak{F}_Y)$ is nontrivial, it is generated by the nontrivial F_i^* and then apply Proposition 3.6. \square

The case when $\text{DG}(\mathfrak{F}_Y)$ is trivial can be recognized by its fixed points.

Definition 3.8 (Fixed points). If $\mathbf{y} \in \text{Per}(\mathfrak{F}_Y)$ is fixed under some simple update order π , then \mathbf{y} must be fixed by each local function F_i . This is because a change to the i^{th} coordinate cannot be corrected by the other local functions in the composition that produces $[\mathfrak{F}_Y, \pi]$. As a consequence, \mathbf{y} is fixed under *all* update orders ω . We write $\text{Fix}(\mathfrak{F}_Y)$ to denote the set of periodic states fixed by some simple update order, or equivalently, the set of periodic states fixed by all local functions F_i (and thus fixed by all update orders).

Proposition 3.9 (Trivial groups and fixed points). *An ω -independent SDS \mathfrak{F}_Y has a trivial dynamics group $\text{DG}(\mathfrak{F}_Y)$ iff $\text{Fix}(\mathfrak{F}_Y) = \text{Per}(\mathfrak{F}_Y)$.*

Proof. If $\text{Fix}(\mathfrak{F}_Y) = \text{Per}(\mathfrak{F}_Y)$ then each F_i^* is trivial and $\text{DG}(\mathfrak{F}_Y)$ is trivial. On the other hand, if $\text{Fix}(\mathfrak{F}_Y) \neq \text{Per}(\mathfrak{F}_Y)$ then there is a periodic state \mathbf{y} and a local function F_i such that $F_i(\mathbf{y}) \neq \mathbf{y}$. For this i , F_i^* is nontrivial and thus $\text{DG}(\mathfrak{F}_Y)$ is nontrivial. \square

When the dynamics group $\text{DG}(\mathfrak{F}_Y)$ is nontrivial, the graph Y can be used to describe the Coxeter group of which it is a quotient.

Definition 3.10 (Coxeter diagrams). The presentation of a Coxeter group is often summarized in graphical form as follows. Given a Coxeter system (W, S) we construct a graph with vertices indexed by S and an edge labeled m_{ij} connecting vertex i and j whenever $m_{ij} > 2$. The Coxeter presentation of W can be easily reconstructed from this edge-labeled graph called the *Coxeter diagram* of W .

Definition 3.11 (Coxeter diagrams for dynamics groups). If \mathfrak{F}_Y is an ω -independent SDS with only two possible vertex states and a nontrivial dynamics group $\text{DG}(\mathfrak{F}_Y)$, then the Coxeter diagram for the Coxeter group of which $\text{DG}(\mathfrak{F}_Y)$ is a quotient can be obtained from Y in three easy steps. First remove every vertex i (and the edges connected to it) for which F_i^* is trivial. Next, remove the edges between i and j when $m_{ij} = 2$ (or equivalently when F_i^* and F_j^* commute). And finally, add the label m_{ij} to each remaining edge. Finally, we note that because F_i^* and F_j^* can only alter coordinates i and j , the cycles of $F_i^* \circ F_j^*$ have size at most 4. Thus each m_{ij} divides 12, the gcd of possible cycle lengths.

Definition 3.12 (Coxeter label). The Coxeter diagram for the Coxeter group mapping onto the dynamics group of an ω -independent Wolfram rule k is particularly simple because of the symmetry of construction of the $\mathfrak{Wol}_n^{(k)}$. Writing F_i instead of $\mathfrak{Wol}_i^{(k)}$ we see that one F_i^* is nontrivial iff they are all nontrivial and the order of $F_i^* \circ F_{i+1}^*$ is a constant independent of i and n . We call this constant c_k , the *Coxeter label* of Wolfram rule k , and as noted above, the value of c_k must divide 12. When G_k is nontrivial all vertices remain. If $c_k > 2$, all edges remain and are labeled c_k .

The result we need from Coxeter theory is an identification of certain classes of groups.

Remark 3.13 (Small Coxeter labels). If $\mathfrak{Wol}_n^{(k)}$ is ω -independent, G_k is nontrivial, and $c_k = 2$, then G_k is a quotient of the Coxeter group \mathbb{Z}_2^n defined by an edgeless Coxeter diagram. If $\mathfrak{Wol}_n^{(k)}$ is ω -independent, G_k is nontrivial, and $c_k = 3$, then G_k is a quotient of the Coxeter group defined by a circular Coxeter diagram with edges labeled 3. This group is the affine Coxeter group of type \tilde{A}_{n-1} but since it is the only Coxeter group we consider without a pre-existing common name (such as SYM_n), we call this group COX_n . Its structure is well-known. Let $(\mathbb{Z}^n \perp \mathbf{1})$ denote the subset of \mathbb{Z}^n perpendicular to the vector $\mathbf{1} = (1, 1, \dots, 1)$, i.e., the set of vectors whose coordinates sum to 0, and note that these form a subgroup under vector addition. If we let SYM_n act on $(\mathbb{Z}^n \perp \mathbf{1})$ by permuting coordinates in the natural way, then COX_n is isomorphic to the semidirect product $(\mathbb{Z}^n \perp \mathbf{1}) \rtimes \text{SYM}_n$. Geometrically, it is the group of isometries of the Euclidean space \mathbb{R}^n that preserve the sum of the coordinates and send vectors with all integer coordinates to other such vectors.

For the sake of concreteness, we select the following explicit isomorphism between COX_n and $(\mathbb{Z}^n \perp \mathbf{1}) \rtimes \text{SYM}_n$. To avoid confusion we use $\mathbf{x} = (x_1, x_2, \dots, x_n)$ for an element of $(\mathbb{Z}^n \perp \mathbf{1})$ and reserve $\mathbf{y} = (y_1, y_2, \dots, y_n)$ for states. Let s_1, s_2, \dots, s_n , denote the Coxeter generators of COX_n . For each $i < n$ let s_i be the element that switches x_i and x_{i+1} , and let s_n send $\mathbf{x} = (x_1, x_2, \dots, x_{n-1}, x_n)$ to $(x_n - 1, x_2, \dots, x_{n-1}, x_1 + 1)$. The natural map to SYM_n only remembers how the subscripts on the x_i s are permuted and ignores the translational component. An element in the kernel of the map $\text{COX}_n \rightarrow \text{SYM}_n$ is called a *pure translation* since it merely adds to \mathbf{x} a vector in $(\mathbb{Z}^n \perp \mathbf{1})$. These pure translations are generated by the elements T_i in COX_n defined as follows. The element $T_1 = s_1 s_2 \cdots s_{n-1} s_n s_{n-1} \cdots s_3 s_2$ and the other T_i are obtained by consistently shifting the subscripts. In terms of its effect on an element $\mathbf{x} \in (\mathbb{Z}^n \perp \mathbf{1})$, T_1 adds the vector $\langle -1, 0, \dots, 0, 1 \rangle$. More generally, T_i adds 1 to x_{i-1} and subtracts 1 from x_i . The pure translations are generated by these commuting T_i which are nearly independent. The sole nontrivial relation they satisfy is that $T_1 + T_2 + \cdots + T_n$ is trivial.

4. TRIVIAL GROUPS

In this section we discuss ω -independent Wolfram rules with trivial dynamics group. Of the 41 representative Wolfram rules, 26 of them fall into this category. For rule 204 (with tag ----), this is immediate since as its tag indicates no local function ever alters the current state. For the other 25 rules the triviality of G_k is a consequence of Proposition 3.9.

More specifically, several lemmas in [2] established conditions under which all periodic points are fixed. These lemmas are listed below along with the representative rules they cover.

- Lemma 5.3 (Rules 0, 4, 8, 12, 72, 76, 128, 132, 136, 140 and 200)
- Lemma 5.5 (Rules 160, 164, 168, 172 and 232)
- Lemma 5.6 (Rules 5, 13, 77, 133 and 141)
- Lemma 6.1 (Rules 32 and 40)
- Lemma 6.3 (Rules 152 and 184)

Even though the group G_k is trivial in each case, the set P_k remains to be calculated. (Many of the proofs in [2] are nonconstructive and do not determine the set of periodic states explicitly.) Because periodic states and fixed states coincide for these rules, there is a relatively straightforward procedure for finding them: simply look at the definition of the rule and remove all states containing triples $y_{i-1}y_iy_{i+1}$ of consecutive states that would lead to an alteration. Clearly the removed states are not fixed by all local functions and do not belong to P_k , and any states remaining at the end of this procedure are fixed by every local function and do belong in P_k . This prompts the following definitions.

Definition 4.1 (Avoiding words). For each n , let $\mathbf{N}_{XY\dots}$ denote the set of states in \mathbb{F}_2^n that do not contain any subwords of the form X, Y, \dots (The letter \mathbf{N} stands for “no”.) For example \mathbf{N}_{11} is the collection of states without adjacent 1s, keeping in mind that we view the subscripts mod n .

Definition 4.2 (Abbreviations). To simplify notation, we introduce six abbreviations: $A='11'$, $B='000'$, $C='111'$, $D='010'$, $E='101'$, and $F='1100'$. Thus, \mathbf{N}_{AE} represents the states in \mathbb{F}_2^n with no subwords of the form ‘11’ or ‘101’. In addition, let $\mathbf{0}$ and $\mathbf{1}$ refer to the state with all 0s and all 1s, respectively.

Remark 4.3 (Why these words). The words we have chosen to abbreviate are those needed to efficiently describe the periodic sets of the 41 representative rules. The words B, C, D and E are triples that need to be avoided, while the words A and F deserve additional explanation. Avoiding the triple ‘110’ is equivalent to avoiding the subword ‘11’ while allowing the state $\mathbf{1}$, and avoiding the triple ‘011’ leads to the same conditions. It thus makes sense to abbreviate the word ‘11’ and treat the state $\mathbf{1}$ separately. The word $F = '1100'$ is only needed to describe P_{28} and P_{29} , so we postpone our discussion of this abbreviation until Section 6. We note that this is the only abbreviated word that is not left-right symmetric.

The periodic sets for the 26 representative Wolfram rules under discussion, calculated as described above, are listed in Table 1. The tag of rule k , along with the decimal of its inversion, reflection and inversion-reflection are also included. Finally we turn to a calculation of the number of orbits for each of these rules.

Remark 4.4 (Recurrence relations). Because the dynamics groups are trivial, we have $o_k = |P_k|$ in each case. Moreover, since each of our sets is defined by a finite list of configurations that it avoids, it is well-known that the number a_n of acceptable configurations

rule	tag	inv	refl	i.r.	o_k	P_k	c_k	G_k
0	0000	255	0	255	1	$\mathbf{0}$	1	1
8	00-0	239	64	253	1	$\mathbf{0}$	1	1
32	x000	251	32	251	1	$\mathbf{0}$	1	1
40	x0-0	235	96	249	1	$\mathbf{0}$	1	1
128	-000	254	128	254	2	$\mathbf{0} \cup \mathbf{1}$	1	1
136	-0-0	238	192	252	2	$\mathbf{0} \cup \mathbf{1}$	1	1
160	1000	250	160	250	2	$\mathbf{0} \cup \mathbf{1}$	1	1
168	10-0	234	224	248	2	$\mathbf{0} \cup \mathbf{1}$	1	1
152	-x-0	230	194	188	2	$\mathbf{0} \cup \mathbf{1}$	1	1
184	1x-0	226	226	184	2	$\mathbf{0} \cup \mathbf{1}$	1	1
4	000-	223	4	223	$ \mathbf{N}_A $	\mathbf{N}_A	1	1
12	00--	207	68	221	$ \mathbf{N}_A $	\mathbf{N}_A	1	1
132	-00-	222	132	222	$ \mathbf{N}_A + 1$	$\mathbf{N}_A \cup \mathbf{1}$	1	1
140	-0--	206	196	220	$ \mathbf{N}_A + 1$	$\mathbf{N}_A \cup \mathbf{1}$	1	1
5	0001	95	5	95	$ \mathbf{N}_{AB} $	\mathbf{N}_{AB}	1	1
13	00-1	79	69	93	$ \mathbf{N}_{AB} $	\mathbf{N}_{AB}	1	1
133	-001	94	133	94	$ \mathbf{N}_{AB} + 1$	$\mathbf{N}_{AB} \cup \mathbf{1}$	1	1
141	-0-1	78	197	92	$ \mathbf{N}_{AB} + 1$	$\mathbf{N}_{AB} \cup \mathbf{1}$	1	1
164	100-	218	164	218	$ \mathbf{N}_{AE} + 1$	$\mathbf{N}_{AE} \cup \mathbf{1}$	1	1
172	10--	202	228	216	$ \mathbf{N}_{AE} + 1$	$\mathbf{N}_{AE} \cup \mathbf{1}$	1	1
77	0--1	77	77	77	$ \mathbf{N}_{BC} $	\mathbf{N}_{BC}	1	1
76	0---	205	76	205	$ \mathbf{N}_C $	\mathbf{N}_C	1	1
72	0--0	237	72	237	$ \mathbf{N}_{CD} $	\mathbf{N}_{CD}	1	1
200	---0	236	200	236	$ \mathbf{N}_D $	\mathbf{N}_D	1	1
232	1--0	232	232	232	$ \mathbf{N}_{DE} $	\mathbf{N}_{DE}	1	1
204	----	204	204	204	2^n	\mathbb{F}_2^n	1	1

TABLE 1. Rules with trivial dynamics.

for each n are the coefficients of an easily calculated rational generating function [5]. As a consequence, they satisfy a constant coefficient recurrence relation. The last statement also follows from [3, Theorem 5.3, p. 132].

In all but two cases these negative descriptions can be reformulated as positive ones that allow us to compute the recurrence with ease. For example, the states in \mathbf{N}_A are those without adjacent 1s and every such state can be uniquely decomposed into subwords of the form ‘0’ and ‘10’ that we call “tiles”. Counting these involve the Lucas numbers and, indirectly, the Fibonacci numbers.

Definition 4.5 (Fibonacci and Lucas). In order to count the size of \mathbf{N}_A we first consider the number b_n of ways to build a word of length n out of the tiles 0 and 10, i.e. with no cyclic subscripting. By focusing on the type of the final tile, we see that the number of such tiles

name	words to avoid	tiles	recurrence relation	Sloane
\mathbf{N}_A	'11'	'0', '10'	$a_n = a_{n-1} + a_{n-2}$	A000032
\mathbf{N}_{AB}	'11', '000'	'10', '100'	$a_n = a_{n-2} + a_{n-3}$	A001608
\mathbf{N}_{AE}	'11', '101'	'0', '100'	$a_n = a_{n-1} + a_{n-3}$	A001609
\mathbf{N}_{BCF}	'000', '111', '1100'	'10', '100', '110'	$a_n = a_{n-2} + 2a_{n-3}$	A072328
\mathbf{N}_{BC}	'000', '111'	'10', '100', '110', '1100'	$a_n = a_{n-2} + 2a_{n-3} + a_{n-4}$	A007040
\mathbf{N}_C	'111'	'0', '10', '110'	$a_n = a_{n-1} + a_{n-2} + a_{n-3}$	A001644
\mathbf{N}_{CD}	'111', '010'	'0', '110'	$a_n = a_{n-1} + a_{n-3}$	A001609
\mathbf{N}_D	'010'	—	$a_n = 2a_{n-1} - a_{n-2} + a_{n-3}$	A109377
\mathbf{N}_{DE}	'010', '101'	—	$a_n = 2a_{n-1} - a_{n-2} + a_{n-4}$	A007039

TABLE 2. Recurrence relations for the basic periodic sets.

satisfy the recurrence relation $b_n = b_{n-1} + b_{n-2}$ with initial conditions $b_1 = 1$ and $b_2 = 2$. The unique solution of the recurrence is $b_n = \text{FIB}_{n+1}$ where FIB_n are the famous *Fibonacci numbers* with values $\{1, 1, 2, 3, 5, 8, 13, 21, \dots\}$ starting with FIB_1 . Returning to the cyclic version, there are three ways the vertex i can be covered by a tile: it can be a 0 tile, the first digit of a 10 tile or the second digit of a 10 tile. Once the tile containing vertex i has been placed, the remaining problem involves tiling a word. Thus $a_n = b_{n-1} + 2b_{n-2}$. It is now easy to see that a_n satisfies the same recurrence relation as b_n but with different initial conditions. In this case $a_1 = 1$ and $a_2 = 3$ and the solution of the recurrence is $a_n = \text{LUC}_n$ where LUC_n are the nearly as famous *Lucas numbers* with values $\{1, 3, 4, 7, 11, 18, 29, \dots\}$ starting with LUC_1 .

This analysis of the size of \mathbf{N}_A easily extends to those cases with a tiling description supplementing the description by words to avoid. The tiles in each case and the corresponding recurrence relation are listed in Table 2. The final column is a reference to the appropriate entry in Neil Sloane's Online Encyclopedia of Integer Sequences [4]. The remaining two cases, \mathbf{N}_D and \mathbf{N}_{DE} , are classical situations where isolated 1s and/or isolated 0s are to be avoided. See the references listed in their entries in [4] for further details.

5. INVERTIBLE RULES

We now turn our attention to Wolfram rules where every state is periodic. Of the 41 representative Wolfram rules, 9 of them fall into this category. Our results for these 9 rules are summarized in Table 3. Note that the more complicated numbers and groups o_k and G_k are not explicitly listed in the table and only described in the text. As a convention, when it is clear that rule k is the rule under discussion, we use F_i instead of the more cumbersome $\text{Wolf}_i^{(k)}$ to denote the rule that updates the i^{th} coordinate. Finally, recall that the Coxeter label c_k must be a divisor of 12 and note that all six possible values occur among these "invertible" rules.

Definition 5.1 (Invertible rules). An SDS \mathfrak{F}_Y is called *invertible* when every state \mathbb{F}^n is periodic, or equivalently when every update rule F_i is a bijection. To see the equivalence

rule	tag	inv	refl	i.r.	o_k	P_k	c_k	G_k
204	----	204	204	204	2^n	\mathbb{F}_2^n	1	1
51	xxxx	51	51	51	1	\mathbb{F}_2^n	2	\mathbb{Z}_2^n
60	xx--	195	102	153	2	\mathbb{F}_2^n	4	$\text{SL}_n(\mathbb{Z}_2)$
150	-xx-	150	150	150	$\lfloor \frac{n}{2} \rfloor + 2$	\mathbb{F}_2^n	3	Thm 5.4
105	x--x	105	105	105	o_{105}	\mathbb{F}_2^n	3	Thm 5.5
156	-x--	198	198	156	$ \mathbf{N}_A \cup \mathbf{1} $	\mathbb{F}_2^n	6	Thm 5.6
201	---x	108	201	108	$ \mathbf{N}_D $	\mathbb{F}_2^n	6	Thm 5.9
57	xx-x	99	99	57	1	\mathbb{F}_2^n	12	Conj 5.10
54	xxx-	147	54	147	2	\mathbb{F}_2^n	12	Conj 5.11

TABLE 3. Invertible rules.

note that when every update rule is bijective, their composition is a permutation and every state is periodic. Conversely, if every state is periodic, then the SDS maps are permutations and the only way this can happen is when every update rule is at least injective and hence bijective.

It is also worth noting that when there are only two possible vertex states, invertibility is characterized by the absence of the symbols 0s and 1s in the tags of the update rules. Before discussing the invertible rules individually, we need some additional notation.

Definition 5.2 (Blocks). If \mathbf{y} is any state other than $\mathbf{0}$ or $\mathbf{1}$ then it consists of alternating strings of 0s and 1s that we abbreviate using exponents. For example $1^5 0^3$ is shorthand for the word 11111000. We call a maximal subword of the form 0^i , 1^j , $0^i 1^j$ or $1^j 0^i$ (with i and j positive) a *0-block*, a *1-block*, a *01-block* and a *10-block*, respectively. For example, keeping in mind the cyclic nature of the subscripts, the state $\mathbf{y} = 1101001$ has two 1-blocks (1 and 1^3), two 0-blocks (0 and 0^2), two 01-blocks (01 and $0^2 1^3$), and two 10-blocks (10^2 and $1^3 0$). Every block has a length. Blocks of length 1 are called *isolated* and longer blocks are called *nontrivial*.

The effect of each of the four possible \mathbf{x} 's in a tag can be described in this language.

Remark 5.3 (Tags and dynamics). When $t_0 = \mathbf{x}$ ($010 \leftrightarrow 000$) isolated 1s can be created and removed and when $t_3 = \mathbf{x}$ ($111 \leftrightarrow 101$) isolated 0s can be created and removed. When $t_1 = \mathbf{x}$ ($011 \leftrightarrow 001$) the boundary between a 1-block and the 0-block to its left can be shifted left or right and when $t_2 = \mathbf{x}$ ($110 \leftrightarrow 100$) the boundary between a 1-block and the 0-block to its right can be shifted left or right.

We now discuss the invertible rules one at a time.

Rule 204. As we already noted in the last section, under rule 204, with tag ----, all of the local functions leave the state of the system unchanged, every local function induces the trivial permutation, and the group generated is the trivial group. There are thus 2^n orbits since every distinct state is an orbit.

Rule 51. Under rule 51, with tag `xxxx`, the local functions ignore their context and always alter the value. Since the local functions pairwise commute $c_{51} = 2$ and the dynamics group is a quotient of \mathbb{Z}_2^n . But since the composition of every distinct subset of local functions toggles a distinct subset of vertex states, there are at least 2^n elements in G_{51} . Thus G_{51} is isomorphic to \mathbb{Z}_2^n . Finally, it is easy to see that there is only one orbit, so $o_{51} = 1$.

Rule 60. The dynamics group G_{60} is interesting because its structure is slightly unexpected. The key observation is that when vertex i is updated, its new value is its old value plus the value of the vertex immediately to its left. In other words $\text{Wolf}_i^{(60)}$ replaces y_i with $y_i + y_{i-1}$, which leads to a matrix representation of the update rules. Viewing \mathbf{y} as a column vector, the effect of updating vertex i can be achieved by multiplying \mathbf{y} on the left by the matrix $\mathbf{A}_i := I + E_{i,i-1}$ where I is the $n \times n$ identity matrix and $E_{i,j}$ is the elementary matrix with 0s everywhere except for a single 1 in the (i, j) position. Matrix multiplication by \mathbf{A}_i is a concise description of the function $\text{Wolf}_i^{(60)}$ from \mathbb{F}_2^n to \mathbb{F}_2^n and thus the matrix group generated by the \mathbf{A}_i s is isomorphic to the dynamics group G_{60} . Since each \mathbf{A}_i has determinant 1, it is clear that G_{60} is a subgroup of $\text{SL}_n(\mathbb{Z}_2)$, and, in fact, it is well-known that these matrices generate all of $\text{SL}_n(\mathbb{Z}_2)$ [6, pg. 455]. Thus G_{60} is isomorphic to $\text{SL}_n(\mathbb{Z}_2)$. Under this group action it is clear that $\mathbf{0}$ is fixed and the remaining states form a single orbit. A calculation shows $c_{60} = 4$.

Rule 150. Rule 150 is similar to rule 60 but this time, when vertex i is updated, its new value is the sum of its old value plus the value of the vertices immediately to its left and its right. This is better known as the *parity function*. More explicitly, $F_i = \text{Wolf}_i^{(150)}$ replaces y_i with $y_{i-1} + y_i + y_{i+1} \pmod 2$, which leads to a matrix representation of the update rules as before. If we define the matrices \mathbf{A}_i as $I + E_{i,i-1} + E_{i,i+1}$ then the group G_{150} can be identified as the subgroup of $\text{SL}_n(\mathbb{Z}_2)$ the matrices \mathbf{A}_i generate. Unlike rule 60 it is not clear which subgroup of $\text{SL}_n(\mathbb{Z}_2)$ this generates. An approach via Coxeter groups is more successful.

Since an easy calculation shows that $c_k = 3$, G_{150} is a quotient of COX_n . Recall from Remark 3.13 that $\text{COX}_n \cong (\mathbb{Z}^n \perp \mathbf{1}) \rtimes \text{SYM}_n$ as well as the conventions established there. The analysis of G_{150} involves two steps. The first is to show that the map $\text{COX}_n \rightarrow G_{150}$ is a factor of the map $\text{COX}_n \rightarrow \text{SYM}_n$ (i.e. the later map decomposes as $\text{COX}_n \rightarrow G_{150} \rightarrow \text{SYM}_n$). To see that G_{150} maps onto SYM_n in a manner consistent with the projection $\text{COX}_n \rightarrow \text{SYM}_n$ start with a state \mathbf{y} and imagine that the numbers 1 up to n are placed in the gaps between the n positions. In particular, initially place the number i between y_{i-1} and y_i . When F_i is applied to a state \mathbf{y} switch the numbers on either side of y_i fixing all the others, in addition to updating the value of y_i . We claim that if i was originally in a gap that marked the end of a 0-block or 1-block, then the same is true of the place where i ends up in the final state. This follows easily from the way rule 150 updates states. From this it is not too hard to see that if a sequence of update rules fixes every state, then each of the numbers 1 through n must also return to their original position. More concretely, if the numbers 1 through n do not all return to their original positions, it is easy to find an explicit state that is not fixed by this sequence of update rules. This means

that there is a well-defined group homomorphism G_{150} to SYM_n that sends the permutation $[\mathfrak{F}_Y, \omega] = F_{\omega_m} \circ \cdots \circ F_{\omega_1}$ to its permutation of the set $\{1, 2, \dots, n\}$. Since this matches the image of $s_{\omega_m} \circ \cdots \circ s_{\omega_1}$ under the map $\text{COX}_n \rightarrow \text{SYM}_n$ we have the factorization we desire.

The existence of maps $\text{COX}_n \rightarrow G_{150} \rightarrow \text{SYM}_n$ imply that the kernel of the first map consists solely of pure translations and we only need to analyze which pure translations lie in the kernel in order to completely understand the group G_{150} . To do this we use the concrete description of COX_n given in Remark 3.13. One of the generating pure translations in COX_n is the element $T_1 = s_1 s_2 \cdots s_{n-1} s_n s_{n-1} \cdots s_3 s_2$. The image of this inside G_{150} is the element $F_1 \circ F_2 \circ \cdots \circ F_{n-1} \circ F_n \circ F_{n-1} \cdots \circ F_3 \circ F_2$. If we apply this sequence of update rules to an arbitrary state \mathbf{y} (using the parity function as we should), the final result is $\mathbf{y} + (y_2 + y_n)\mathbf{1}$. More generally, the image of T_i acts on states by sending \mathbf{y} to $\mathbf{y} + (y_{i-1} + y_{i+1})\mathbf{1}$. Applying T_i twice is clearly trivial so the kernel contains the subgroup of $(\mathbb{Z}^n \perp \mathbf{1})$ that the vectors $2T_i$ generate. This set is $(2\mathbb{Z})^n \perp \mathbf{1}$. As a consequence, G_{150} is a quotient of the group $(\mathbb{Z}_2^n \perp \mathbf{1}) \rtimes \text{SYM}_n$ of size $2^{n-1} \cdot n!$. When n is even there is another pure translation in the kernel, namely, the result of applying once each T_i with an odd subscript. As a pure translation this adds the vector $\langle -1, 1, -1, \dots, -1, 1 \rangle$. Equivalently, when n is even G_{150} is a quotient of $((\mathbb{Z}_2^n \perp \mathbf{1})/\langle \mathbf{1} \rangle) \rtimes \text{SYM}_n$. (Note that adding $\mathbf{1}$ to a state \mathbf{x} replaces \mathbf{x} with its *complement*, the 0s become 1s and vice versa. Moreover, this pure translation commutes with the symmetric group action and is central in $(\mathbb{Z}_2^n \perp \mathbf{1}) \rtimes \text{SYM}_n$.) We now show that these are the only pure translations in the kernel.

Theorem 5.4 (Rule 150). *When n is odd the group G_{150} is isomorphic to $(\mathbb{Z}_2^n \perp \mathbf{1}) \rtimes \text{SYM}_n$ and when n is even it is isomorphic to $((\mathbb{Z}_2^n \perp \mathbf{1})/\langle \mathbf{1} \rangle) \rtimes \text{SYM}_n$. In particular, $|G_{150}| = 2^{n-1} \cdot n!$ when $2 \nmid n$ and $2^{n-2} \cdot n!$ when $2 \mid n$.*

Proof. That G_{150} is a quotient of these groups was shown above, so we only need to show that we have found the full kernel of the map $\text{COX}_n \rightarrow G_{150}$. To see whether there are any other pure translations that are trivial in G_{150} note that T_{i-1} and T_{i+1} are the only generating pure translations that add $y_i \mathbf{1}$ to \mathbf{y} . Thus, if $a_1 T_1 + a_2 T_2 + \cdots + a_n T_n$ acts trivially on every state \mathbf{y} , the parity of a_{i-1} and a_{i+1} must match for every i . When n is odd, this mean that all the a_i have the same parity and every such element is one we already know lies in the kernel. When n is even there is one additional possibility. Perhaps the a_i s with even subscripts have one parity and the ones with odd subscripts have the other. Removing summands we already know to lie in the kernel and using the relation $T_1 + T_2 + \cdots + T_n = \mathbf{0}$ if necessary, we see that this possibility is equivalent to $T_1 + T_3 + T_5 + \cdots + T_{n-1}$, again, an element we already know lies in the kernel. \square

Although we did not need its orbit structure in order to analyze G_{150} , it is easy to see that the changes rule 150 allows (and the only changes it allows) are the alterations of the boundaries of the 0-blocks and 1-blocks. Thus two states belong to the same orbit iff they have the same number of 0-blocks and the same number of 1-blocks. In particular $\alpha_{150} = \lfloor \frac{n}{2} \rfloor + 2$ with the 2 corresponding to the fixed states $\mathbf{0}$ and $\mathbf{1}$.

Rule 105. Rule 105 is the negation of the parity function and its analysis is very similar to our analysis of rule 150. In particular, $F_i = \text{Wolf}_i^{(105)}$ replaces y_i with $1 + y_{i-1} + y_i + y_{i+1}$

mod 2. The group G_{105} can be described as a subgroup of a linear group (with an extra row and column added for constants) but viewing it as a Coxeter quotient is more fruitful. The Coxeter label $c_{105} = 3$, and $\text{COX}_n \rightarrow G_{105} \rightarrow \text{SYM}_n$ as above. The argument for such a factorization is similar in spirit to the one presented above for rule 150 but complicated by the presence of negations. We omit the details.

The existence of maps $\text{COX}_n \rightarrow G_{105} \rightarrow \text{SYM}_n$ imply that the kernel of the first map consists solely of pure translations and we only need to analyze which pure translations lie in the kernel in order to completely understand the group G_{105} . This time the image of the pure translation T_i in G_{105} acts on states by sending \mathbf{y} to $\mathbf{y} + (1 + y_{i-1} + y_{i+1})\mathbf{1}$. Applying T_i twice is clearly trivial so the kernel contains the subgroup of $(\mathbb{Z}^n \perp \mathbf{1})$ that the vectors $2T_i$ generate. This set is $(2\mathbb{Z})^n \perp \mathbf{1}$. As a consequence, G_{105} is a quotient of the group $(\mathbb{Z}_2^n \perp \mathbf{1}) \rtimes \text{SYM}_n$ of size $2^{n-1} \cdot n!$. When n is a multiple of 4 there is another pure translation in the kernel, namely, the result of applying once each T_i with an odd subscript. The reason n needs to be a multiple of 4 and not merely even is that when n is twice an odd number \mathbf{y} is sent to $\mathbf{y} + \mathbf{1}$. As a pure translation this adds the vector $\langle -1, 1, -1, \dots, -1, 1 \rangle$. Equivalently, when n is even G_{105} is a quotient of $((\mathbb{Z}_2^n \perp \mathbf{1})/\langle \mathbf{1} \rangle) \rtimes \text{SYM}_n$. We now show that these are the only pure translations in the kernel.

Theorem 5.5 (Rule 105). *When n is not a multiple of 4 the group G_{105} is isomorphic to $(\mathbb{Z}_2^n \perp \mathbf{1}) \rtimes \text{SYM}_n$ and when it is a multiple of 4 it is isomorphic to $((\mathbb{Z}_2^n \perp \mathbf{1})/\langle \mathbf{1} \rangle) \rtimes \text{SYM}_n$. In particular, $|G_{105}| = 2^{n-1} \cdot n!$ when $4 \nmid n$ and $|G_{105}| = 2^{n-2} \cdot n!$ when $4 \mid n$.*

Proof. That G_{105} is a quotient of these groups was shown above, so we only need to show that we have found the full kernel of the map $\text{COX}_n \rightarrow G_{105}$. To see whether there are any other pure translations that are trivial in G_{105} note that T_{i-1} and T_{i+1} are the only generating pure translations that add $y_i\mathbf{1}$ to \mathbf{y} . Thus, if $a_1T_1 + a_2T_2 + \dots + a_nT_n$ acts trivially on every state \mathbf{y} , the parity of a_{i-1} and a_{i+1} must match for every i . When n is odd, this mean that all the a_i have the same parity and every such element is one we already know lies in the kernel. When n is even there is one additional possibility. Perhaps the a_i s with even subscripts have one parity and the ones with odd subscripts have the other. Removing summands we already know to lie in the kernel and using the relation $T_1 + T_2 + \dots + T_n = \mathbf{0}$ if necessary, we see that this possibility is equivalent to $T_1 + T_3 + T_5 + \dots + T_{n-1}$. If n is twice an odd number then \mathbf{y} is sent to $\mathbf{y} + \mathbf{1}$ and this is not in the kernel, but when n is twice an even number (i.e. $4 \mid n$) then this is an element we already know lies in the kernel. \square

Finally, for the sake of completeness, we include a description of o_{105} but leave a verification of our assertions as an exercise. When n is twice an odd number $o_{105} = 2 \lfloor \frac{n}{4} \rfloor + 4$, when n is twice an even number $o_{105} = 2 \lfloor \frac{n}{4} \rfloor + 2$, and when n is odd, $o_{105} = 2$ (and these are the orbits of $\mathbf{0}$ and $\mathbf{1}$).

Rule 156. Wolfram rule 156 is a case where the orbit structure is useful and we use it to prove that G_{156} is a subdirect product of symmetric groups of particular sizes in a very concrete fashion. The first thing to notice is that the only changes rule 156 allows are growing or shrinking 1-blocks from the right. Isolated 1s and 0s can neither be created nor

removed (so the number of blocks is an invariant) and the left end of a 1-block is never moved (so that subwords of the form 01 persist forever). To analyze the number of orbits, apply the rewrite rule $110 \rightarrow 100$ to remove as many 1s as possible. It should be clear that (unless we started with the state $\mathbf{1}$) this process ends when all remaining 1s are isolated. Moreover, these remaining 1 are precisely the leftmost 1s in the initial blocks so there is a *unique* state in each orbit with isolated ones. In other words, the orbits are in natural bijection with the set $\mathbf{N}_A \cup \mathbf{1}$, and thus $o_{156} = |\mathbf{N}_A \cup \mathbf{1}| = \text{LUC}_n + 1$.

We now examine the way the updates rules act on the states in each orbit individually. The fixed state $\mathbf{1}$ can be discarded since the trivial group it generates does not contribute meaningfully to the subdirect product. Next we note that the dynamics group restricted to a single orbit usually splits further as a direct product. For example, consider the state $\mathbf{y} = 1000001000$. It has a two 10-blocks, one of length 6 and one of length 4. The orbit of \mathbf{y} consists of all words that can be written as a 10-block of length 6 followed by a 10-block of length 4. In particular, the first six digits can be 100000, 110000, 111000, 111100 or 111110 and the last four can be 1000, 1100 or 1110. These fifteen combinations form the complete orbit. In addition, writing F_i in place of $\text{Wolf}_i^{(156)}$, the only update rules that act nontrivially on this orbit are F_2, F_3, F_4, F_5, F_8 and F_9 , but F_2, F_3, F_4 and F_5 commute with F_8 and F_9 so G_{156} restricted to this orbit splits as a direct product $\langle F_2, F_3, F_4, F_5 \rangle \times \langle F_8, F_9 \rangle$.

If we focus in on the first six digits for a moment and name the possibilities $x_1 = 100000$, $x_2 = 110000$, $x_3 = 111000$, $x_4 = 111100$ and $x_5 = 111110$ then we see that F_2 swaps x_1 and x_2 and fixes the other x_i , F_3 swaps x_2 and x_3 , F_4 swaps x_3 and x_4 and F_5 swaps x_4 and x_5 . The induced permutations of the five possibilities are $F_2 = (12)$, $F_3 = (23)$, $F_4 = (34)$ and $F_5 = (45)$. Thus, these four update rules, restricted in this way generate a copy of SYM_5 with the functions F_2, F_3, F_4 and F_5 acting as the standard Coxeter generating set.

More generally, if a state \mathbf{y} contains a 10-block of length $m + 1$ starting at position ℓ , then every state in the orbit of \mathbf{y} contains a 10-block of this length at this location. There are m possibilities for a block of this type that we call $x_i = 1^i 0^j$ where $i, j > 0$ and $i + j = m + 1$. The update rules $F_{\ell+1}, \dots, F_{\ell+m-1}$ are the only ones that alter this block and they generate a copy of SYM_m acting on the possibilities x_i in the standard way. We call this group $\text{SYM}_m^{(\ell)}$. As a group it is SYM_m and the number ℓ indicates which update rules generate SYM_m and in what way. Returning to our earlier example, we see that the quotient of G_{156} obtained by restricting to the orbit of $\mathbf{y} = 1000001000$ is the group $\text{SYM}_5^{(1)} \times \text{SYM}_3^{(7)}$. More generally, if \mathbf{y} has 10-blocks of lengths $m_1 + 1, m_2 + 1, \dots, m_k + 1$ starting at $\ell_1, \ell_2, \dots, \ell_k$ then the quotient of G_{156} restricted to the orbit of \mathbf{y} is the group $\text{SYM}_{m_1}^{(\ell_1)} \times \text{SYM}_{m_2}^{(\ell_2)} \times \dots \times \text{SYM}_{m_k}^{(\ell_k)}$.

To summarize, G_{156} is a subdirect product of groups, each of which is isomorphic to a direct product of groups of the form $\text{SYM}_m^{(\ell)}$. When written out completely using the orbit structure, the groups $\text{SYM}_m^{(\ell)}$ occur multiple times and the projection of G_{156} onto each of these repeated factors is identical. As a consequence, the group G_{156} can be embedded in a direct product where each of these groups only occurs once. Note that the position ℓ can be any number 1 through n , but that $m + 1$ can only be 2 through n excluding $n - 1$ since

n	$ G_{156} $	$ G_{201} $	$ G_{73} $
4	$2^3 \cdot 3^4$	$(7!/2)$	$(7!/2)$
5	$2^{15} \cdot 3^5$	$2^5 \cdot (11!/2)$	$2^5 \cdot (11!/2)$
6	$2^{18} \cdot 3^{12} \cdot 5^6$	$2^6 \cdot 3^6 \cdot (18!/2)$	$2^5 \cdot 3^6 \cdot (18!/2)$
7	$2^{42} \cdot 3^{28} \cdot 5^7$	$2^{21} \cdot 3^{14} \cdot 5^7 \cdot (29!/2)$	$2^{21} \cdot 3^7 \cdot 5^7 \cdot (29!/2)$

TABLE 4. The orders of G_{156} , G_{201} and G_{73} for small values of n .

the complement of the 10-block with length $m + 1$ must be tilable by 10-blocks and these have length at least 2.

Theorem 5.6 (Rule 156). *The group G_{156} is a subdirect product of symmetric groups. In particular,*

$$G_{156} \subset \prod_{\ell=1}^n \left(\text{SYM}_{n-1}^{(\ell)} \times \prod_{m=2}^{n-3} \text{SYM}_m^{(\ell)} \right)$$

and each generator of G_{156} restricted to a factor is either trivial or a standard Coxeter generator as described above.

Note that the factors with $m = 1$ have been eliminated from these products, which is possible because the groups $\text{SYM}_1^{(\ell)}$ are trivial. We computed the order of G_{156} for $4 \leq n \leq 7$, and the results are shown in Table 4. Notice that for $n > 4$ each of these orders are n^{th} powers.

Rule 201. Rule 201 is a second situation where the orbit structure can be used to simplify our analysis. In particular we use orbit structures to prove that G_{201} is a subdirect product of groups that are all (conjecturally) symmetric or alternating groups of particular sizes in a very concrete fashion. The first step is to analyze the orbit structure. The only changes under rule 201 are the creation or removal of isolated 1s. In particular the size and location of nontrivial 1-blocks in a state \mathbf{y} are an invariant of its orbit. In fact, since we can remove all isolated 1s, any two states with the same set of nontrivial 1-blocks belong to the same orbit and every orbit has exactly one state with no isolated 1s. Thus $o_{201} = |\mathbf{N}_D|$, the number of states with no isolated 1s.

Let \mathbf{y} be a state with no isolated 1s. It should be clear that the restriction of the group G_{201} to the orbit of \mathbf{y} will split as a direct product with one factor for each 0-block in \mathbf{y} . This is because all of the local functions that insert and remove isolated 1s in one 0-block will commute with those that insert and remove isolated 1s in a different 0-block, being separated by nontrivial 1-blocks. As was the case with rule 156, the orbit groups of G_{201} split into factors, the refined direct product contains redundancies and these redundancies can be removed to give a better subdirect product representation of G_{201} . In particular, G_{201} embeds in the subdirect product of orbit groups acting on orbits with only one 0-block. Let LG_n denote the group obtained by restricting G_{201} to the orbit of $\mathbf{0}$ and let $\text{FG}_m^{(\ell)}$ be the group obtained by restricting G_{201} to the orbit of the state \mathbf{y} with a single 0-block of length m starting at position ℓ . The notations stand for *Lucas group* and *Fibonacci group*

and are suggested by the fact that the orbit of $\mathbf{0}$ under G_{201} is \mathbf{N}_A of size LUC_n and the orbit of \mathbf{y} under G_{201} has size FIB_m (the first 0 is fixed and the remaining $m - 1$ digits of the 0-block are tiled by tiles 0 and 10). In the case of LG_n we can be slightly more precise.

Proposition 5.7 (Lucas groups). *The group LG_n is a subgroup of $\text{SYM}_{\text{LUC}_n}$ when FIB_{n-1} is odd and a subgroup of $\text{ALT}_{\text{LUC}_n}$ when FIB_{n-1} is even.*

Proof. Since LG_n acts on a set of size $|\mathbf{N}_A| = \text{LUC}_n$, only the second assertion needs to be established. Let F_i^* denote the restriction of $\text{Wolf}_i^{(201)}$ to \mathbf{N}_A and note that these permutations generate LG_n . To count how many 2-cycles are in the disjoint cycle notation for F_i^* , note that for each $\mathbf{y} \in \mathbf{N}_A$, $F_i^*(\mathbf{y}) \neq \mathbf{y}$ iff $y_{i-1} = y_{i+1} = 0$. Thus, the number of 2-cycles in F_i^* equals the number of words $y_{i+1}y_{i+2} \cdots y_{i-2}y_{i-1}$ of length $n - 1$ with only isolated 1s and $y_{i+1} = y_{i-1} = 0$, which equals the number of ways to tile a word $y_{i+2} \cdots y_{i-2}y_{i-1}$ of length $n - 2$ with tiles 0 and 10. As we saw in Definition 4.5, this number is FIB_{n-1} . Thus F_i^* is an even permutation iff FIB_{n-1} is even, and when this is true, LG_n is a subgroup of $\text{ALT}_{\text{LUC}(n)}$. \square

The following conjecture is based on computational evidence for small values of n . We have checked, for example, that $\text{FG}_6^{(\ell)} \cong \text{SYM}_8 = \text{SYM}_{\text{FIB}_6}$ and that $\text{LG}_4 \cong \text{ALT}_7 = \text{ALT}_{\text{LUC}_4}$ (as predicted since $\text{FIB}_3 = 2$ is even).

Conjecture 5.8 (Fibonacci groups and Lucas groups). *For all m , $\text{FG}_m^{(\ell)} \cong \text{SYM}_{\text{FIB}_m}$ and for all n , $\text{LG}_n \cong \text{SYM}_{\text{LUC}_n}$ when FIB_{n-1} is odd and $\text{LG}_n \cong \text{ALT}_{\text{LUC}_n}$ when FIB_{n-1} is even.*

The parity of the Fibonacci numbers is quite predictable and a more direct statement is possible, but the one given is more closely tied to the reason LG_n stays inside $\text{ALT}_{\text{LUC}_n}$. The obstruction to establishing Conjecture 5.8 is how far the permutations generating the Fibonacci groups and Lucas groups differ from the standard generating sets of the symmetric and alternating groups. In any case, this analysis of the orbit structure under rule 201 gives the following subdirect product decomposition for G_{201} . The values of m only range from 3 to $n - 2$ to exclude trivial groups on the low end and because the complement of the 0-block needs to leave room for a nontrivial block of 1s.

Theorem 5.9 (Rule 201). *The group G_{201} is a subdirect product of Lucas groups and Fibonacci groups. In particular,*

$$G_{201} \subset \text{LG}_n \times \prod_{\ell=1}^n \prod_{m=3}^{n-2} \text{FG}_m^{(\ell)}$$

Finally, we used a computer program to calculate the order of G_{201} for $4 \leq n \leq 7$, and the results are shown in Table 4. Notice that the answer in each case is an n^{th} power times the size of $\text{ALT}_{\text{LUC}_n}$.

Rule 57. Wolfram rule 57 can introduce and remove isolated 0s and 1s and it can grow and shrink 0-blocks from the left and 1-blocks from the right. With so much flexibility it is easy to see that there is only one orbit and thus $o_{57} = 1$. Moreover, because $\text{Wolf}_i^{(57)}(\mathbf{y}) = \mathbf{y}$ only when $y_{i-1} = 0$ and $y_{i+1} = 1$, the number of 2-cycles in the disjoint cycle representation

of $\text{Wolf}_i^{(57)}$ is 2^{n-3} . In particular, it is an even permutation for $n > 4$ and we conclude that G_{57} is a subgroup of ALT_{2^n} .

Conjecture 5.10 (Rule 57). *For all $n > 4$, $G_{57} \cong \text{ALT}_{2^n}$.*

We have verified Conjecture 5.10 explicitly up to $n = 8$. The reason Conjecture 5.10 is difficult to establish abstractly is that the permutations $\text{Wolf}_i^{(57)}$ are far removed from the standard generating sets of ALT_{2^n} .

Rule 54. Rule 54 is similar to rule 57 except that the state $\mathbf{0}$ is now fixed. Since 1-blocks can grow to the left or the right and isolated 0s can be removed, all states other than $\mathbf{0}$ are in the same orbit as $\mathbf{1}$. Thus $o_{54} = 2$. Since fixed states only contribute trivial groups to the subdirect product structure, they can be ignored when computing the dynamics group. In particular, G_{54} is a subgroup of $\text{SYM}_{2^{n-1}}$. Because $\text{Wolf}_i^{(54)}(\mathbf{y}) = \mathbf{y}$ only when $y_{i-1} = y_{i+1} = 0$, the number of 2-cycles in the disjoint cycle representation of $\text{Wolf}_i^{(54)}$ is once again 2^{n-3} , which is even for $n > 4$, and G_{54} lies in $\text{ALT}_{2^{n-1}}$.

Conjecture 5.11 (Rule 54). *For all $n > 4$, $G_{54} \cong \text{ALT}_{2^{n-1}}$.*

We have verified Conjecture 5.11 explicitly up to $n = 8$ as well. The reason Conjecture 5.11 is difficult to establish abstractly is that, as with rule 57, the permutations $\text{Wolf}_i^{(54)}$ are far removed from the standard generating sets of $\text{ALT}_{2^{n-1}}$.

6. SEVEN EXCEPTIONS

Of the 41 representative Wolfram rules, we have seen that 26 have trivial dynamics groups, 9 are invertible, and 1 rule is both. In this section we study the remaining 7 rules. One key to understanding the dynamics of these noninvertible rules with nontrivial dynamics is to note that they agree with an invertible rule on a union of its orbits, which implies that the dynamics group of the noninvertible rule is a homomorphic image of the dynamics group of the invertible rule with which it agrees. We begin by establishing these general facts.

Theorem 6.1 (Independence and invertibility). *If \mathfrak{F}_Y is an ω -independent SDS, then there exists an invertible SDS \mathfrak{F}'_Y such that (1) \mathfrak{F}_Y and \mathfrak{F}'_Y agree on the set $\text{Per}(\mathfrak{F}_Y)$, (2) $\text{Per}(\mathfrak{F}_Y)$ is a union of orbits of the action of $\text{DG}(\mathfrak{F}'_Y)$ on \mathbb{F}^n , and (3) there is a surjection from $\text{DG}(\mathfrak{F}'_Y)$ onto $\text{DG}(\mathfrak{F}_Y)$.*

Proof. We define new local functions F'_i as follows. Call a state $\mathbf{y} \in \mathbb{F}^n$ *i*-periodic if the sequence $\mathbf{y}, F_i(\mathbf{y}), F_i(F_i(\mathbf{y})), F_i(F_i(F_i(\mathbf{y})))$, etc., eventually returns to \mathbf{y} ; otherwise call it *i*-transitory. The function F'_i is defined as equal to F_i on the *i*-periodic states and the identity function of the *i*-transitory ones. It is now straightforward to check that F'_i is both *Y*-local at *i* and bijective. By Definition 5.1, the bijectivity of each F'_i means that \mathfrak{F}'_Y is invertible. Next, by Proposition 1.7 each F_i^* is a permutation of $\text{Per}(\mathfrak{F}_Y)$. Thus the states in $\text{Per}(\mathfrak{F}_Y)$ are *i*-periodic for each *i*, and \mathfrak{F}_Y and \mathfrak{F}'_Y agree completely on this set of states. The invariance of $\text{Per}(\mathfrak{F}_Y)$ under each F_i and its agreement with the corresponding

Rule	tag	inv	refl	i.r.	o_k	P_k	c_k	G'_k
28	0x--	199	70	157	$ N_{AB} + 1$	$N_{BCF} \cup \mathbf{0}$	2	\mathbb{Z}_2^n
29	0x-1	71	71	29	$ N_{AB} $	N_{BCF}	2	\mathbb{Z}_2^n
1	000x	127	1	127	1	N_A	6	LG_n
9	00-x	111	65	125	1	N_A	6	LG_n
129	-00x	126	129	126	2	$N_A \cup \mathbf{1}$	6	LG_n
137	-0-x	110	193	124	2	$N_A \cup \mathbf{1}$	6	LG_n
73	0--x	109	73	109	$ N_{CD} $	N_C	6	Thm 6.5

TABLE 5. The seven exceptional cases.

F'_i means that $\text{Per}(\mathfrak{F}_Y)$ is a union of orbits under the action of $\text{DG}(\mathfrak{F}'_Y)$. And finally, the projection from the subdirect product containing $\text{DG}(\mathfrak{F}'_Y)$ onto the factors corresponding to the orbits contained in $\text{Per}(\mathfrak{F}_Y)$ produces the required surjective group homomorphism. \square

Corollary 6.2 (Independence and invertibility). *If \mathfrak{F}_Y is an ω -independent SDS with only two possible vertex states, then the update rules for an invertible SDS \mathfrak{F}'_Y satisfying the conclusions of Theorem 6.1 can be obtained by replacing every 0 and 1 in the tags of the update rules of \mathfrak{F}_Y with the symbol -.*

Proof. The assertion merely describes the effect the previous construction has on tags. \square

As an illustration, we apply the corollary to each of the 7 exceptional rules.

Example 6.3 (Exceptional rules and invertibility). By Corollary 6.2, rules 28 (0x--) and 29 (0x-1) agree with rule 156 (-x--) when restricted to P_{28} and P_{29} , respectively. Similarly, rules 1 (000x), 9 (00-x), 129 (-00x), 137 (-0-x) and 73 (0--x) agree with rule 201 (---x) when restricted to P_1 , P_9 , P_{129} , P_{137} , and P_{73} , respectively.

The proof of Theorem 6.1 also provides a concrete description of the periodic states.

Corollary 6.4 (Periodic orbits). *If \mathfrak{F}_Y is an ω -independent SDS with only two possible vertex states and \mathfrak{F}'_Y is the corresponding invertible SDS satisfying the conclusions of Theorem 6.1, then the nonperiodic states under \mathfrak{F}_Y are the union of the orbits under \mathfrak{F}'_Y that contain a state \mathbf{y} that can be altered as a consequence of a 0 or a 1 in the tag of an update rule of \mathfrak{F}_Y .*

Table 5 contains a summary of our results about these 7 exceptional rules. These noninvertible rules with nontrivial dynamics naturally fall into three classes: $\{28, 29\}$, $\{1, 9, 129, 137\}$, and $\{73\}$.

Rules 28 and 29. By Corollary 6.2, both G_{28} and G_{29} are quotients of G_{156} . The first step is to decide which states are in P_{28} and P_{29} . By Corollary 6.4, rule 28 (with tag 0x--) must exclude all orbits that contain a state with subword 111 and rule 29 (with tag 0x-1) must exclude all orbits that contain a state with subword 111 or 000. For rule 29, we are left with those states that have no 10 blocks of length 4 or more (because any orbit that

contains 1000 or 1100 also contains 1110) and for rule 28 we are left with the same states plus $\mathbf{0}$ which is fixed. These states can be described by excluding $B = \text{'111'}$, $C = \text{'000'}$ and $F = \text{'1100'}$. Because fixed states contribute trivial groups to the subdirect product and both G_{28} and G_{29} are obtained by restricting the action of G_{156} on \mathbb{F}_2^n to the same set of non-fixed periodic states, G_{28} and G_{29} are isomorphic groups. The number of orbits depends on the number of ways to tile n cyclically arranged positions by tiles of size 2 and 3 and, as we saw in Section 4, this number is counted by $|N_{AB}|$. Thus $o_{29} = |N_{AB}|$ and $o_{28} = |N_{AB}| + 1$. Continuing the notation used in our discussion of invertible rule 156, this means that $G_{28} = G_{29}$ is the image of G_{156} in the product $\text{SYM}_2^{(1)} \times \text{SYM}_2^{(2)} \times \cdots \times \text{SYM}_2^{(n)}$ which is isomorphic to \mathbb{Z}_2^n . And since G_{156} projects fully onto each factor, $G_{28} = G_{29}$ is all of \mathbb{Z}_2^n . As a final note, when $n = 4$, $m + 1$ is not allowed to be 3 and the factors SYM_2^ℓ do not occur in the subdirect product. Thus $G_{28} = G_{29}$ are trivial groups for $n = 4$.

An alternative way to proceed would have been to calculate $c_{28} = c_{29} = 2$, conclude that $G_{28} = G_{29}$ is a quotient of \mathbb{Z}_2^n and then to argue that it is actually equal to \mathbb{Z}_2^n . Finally, we remark that rules 28 and 29 have the same dynamics group as rule 51 but for very different reasons. It is perhaps surprising that rules 28 and 29 have such simple dynamics groups since they were among the six exceptional cases that needed to be dealt with separately in the proof of Theorem 2.4 due to the complications in analyzing their dynamics.

Rules 1, 9, 129 and 137. By Corollary 6.2, G_1 , G_9 , G_{129} and G_{137} are all quotients of G_{201} . To decide which quotient, the first step is to decide which states are in P_1 , P_9 , P_{129} and P_{137} . By Corollary 6.4, all four sets must exclude any state with adjacent ones (although P_{129} and P_{137} include the fixed state $\mathbf{1}$) and, by our previous analysis of rule 201, the states with only isolated 1s form a single orbit under its action and thus are included in all four periodic sets. This means that $o_1 = o_9 = 1$ and $o_{129} = o_{137} = 2$. Moreover, recalling that fixed states contribute trivial groups to the subdirect product, G_1 , G_9 , G_{129} and G_{137} are all isomorphic to the group obtained when the action of G_{201} is restricted to the single orbit N_A . In other words, the groups G_1 , G_9 , G_{129} , and G_{137} are all isomorphic to LG_n .

Rule 73. Our final rule, rule 73, is also a quotient of G_{201} as indicated by Corollary 6.2, but this time the number of orbits retained is much larger. In fact, by Corollary 6.4, only states containing $C = \text{'111'}$ as a subword need to be removed. Since neither rule 201 nor rule 73 can bring subwords of the form '111' into existence, all states in N_C are periodic under G_{73} . Because the orbits under rule 201 each contain a unique state with no isolated 1s, the number of orbits in P_{73} equals the number of states with no '111' and no '010'. Thus $P_{73} = N_C$ and $o_{73} = |N_{CD}|$. When the orbit groups for these orbits are factored into direct products and redundancies are removed, the Lucas group LG_n and all of the Fibonacci groups $\text{FG}_m^{(\ell)}$ arise, except those with $m = n - 3$ and $m = n - 4$. This is because the complement of the 0-block of length m starting at location ℓ must be tilable with pairs of adjacent 1s that alternate with other 0-blocks. When $m = n - 2$ the complement is a single pair of 1s and when $m \leq n - 5$ the complement can be a word of the form $1^{2^i}1^2$ with $i > 0$.

Theorem 6.5. *The group G_{73} is the projection of G_{201} into the following direct product:*

$$G_{73} \subset \text{LG}_n \times \prod_{\ell=1}^n \left(\text{FG}_{n-2}^{(\ell)} \times \prod_{m=3}^{n-5} \text{FG}_m^{(\ell)} \right)$$

As with rule 201, we used a computer program to calculate the order of G_{73} for $4 \leq n \leq 7$, and the results are shown in Table 4.

7. CONCLUSION

Now that the periodic states and the dynamics groups have been roughly described for every ω -independent ACA, there are two types of generalizations that one might pursue. One natural extension would be to take a closer look at the Wolfram rules that are *dependent* on the update order ω and study the variety of periodic states and dynamics that arise. Because of the dependence on update order, however, dynamics groups are no longer well-defined in this context and alternative concepts will need to be developed. Alternatively, as we have tried to highlight throughout, many of the techniques developed here apply to larger classes of ω -independent SDSs. As in [2], we started with ACAs because they are simple enough to be analyzed but complicated enough to exhibit interesting behavior. Analyzing SDSs over more complicated graphs with a greater number of vertex states pose a bigger challenge.

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