

## Adjoints

Def. We say  $F: \mathcal{C} \rightarrow \mathcal{D}$ ,  $G: \mathcal{D} \rightarrow \mathcal{C}$  are an adjoint pair ( $F$  is left adjoint to  $G$ ) if there is a natural iso

$$\mathcal{D}(fx, y) \cong \mathcal{C}(x, Gy)$$

Functors  $\mathcal{C} \times \mathcal{D}^{\text{op}} \rightarrow \underline{\text{Set}}$ .

Ex Vector spaces

Fix a field  $k$ . Define  $F: \underline{\text{Set}} \rightarrow k\text{-Vect}$  by taking the  $k$ -vector space with basis elements of the set.

$$F\{a, b, c\} = \{ \underset{\text{finite}}{\substack{\uparrow \\ k\text{-linear combos of } a, b, c}} \}$$

$G: k\text{-Vect} \rightarrow \underline{\text{Set}}$  the forgetful functor:  
takes the vector space to itself considered as a set.

$f \in \text{Set}(x, y)$ , then  $Ff: fx \rightarrow fy$  is defined by taking any basis elt in  $fx$  to the basis elt from  $f$ : "linearly extend  $f$ ".

$$\underline{k\text{-Vect}}(FS, V) \cong \underline{\text{Set}}(S, GV)$$

$\nwarrow \text{Set}$

A morphism of vector spaces is determined by where it sends the basis ( $S$ ).  
canon.

$$\begin{array}{ccc} S & \xrightarrow{\quad} & GFS \\ f \searrow & \downarrow \exists! k\text{-linear map} & \\ & GV & \text{extending } f \end{array}$$

$$\underline{\text{Ex: Met}} \quad \underline{\text{CMpt}}$$

$\uparrow \quad \uparrow$

morphisms are uniformly continuous

If I have a uniformly continuous  $f: X \rightarrow Y$ ,  
 complete  $\exists! \tilde{f}: \tilde{X} \rightarrow Y$ , where  $\tilde{X}$   
 is the completion of  $X$ , s.t.

$$\begin{array}{ccc} X & \xrightarrow{\quad} & \tilde{X} \\ f \searrow & \downarrow \tilde{f} & \rightarrow \exists! \\ & Y & \end{array}$$

Completion functor  $F$ , forgetful  $G: \underline{\text{CMet}} \rightarrow \underline{\text{Met}}$

$$\text{Met}(X, GY) \cong \underline{\text{CMet}}(FX, Y)$$

Completion is left adjoint to forgetful.

Lemma. If  $F$  is left adjoint to  $G$  ( $F \dashv G$ ),  
then  $GF : \mathcal{C} \rightarrow \mathcal{C}$ ,  $FG : \mathcal{D} \rightarrow \mathcal{D}$ .

We get a natural transformation  $\eta : 1_{\mathcal{C}} \rightarrow GF$ ,  
 $\epsilon : FG \rightarrow 1_{\mathcal{D}}$ , instantiating the ISO:

$$\begin{array}{ccc} \mathcal{D}(Fx, y) & \longrightarrow & \mathcal{C}(x, Gy) \\ f \downarrow & \longrightarrow & Gf \circ \eta_x \end{array}$$

$$\begin{array}{ccc} x & , & Gf \circ \eta_x \\ \eta_x \downarrow & & \downarrow \\ GFx & \xrightarrow{Gf} & Gy \end{array}$$

$$\begin{array}{ccc} \mathcal{C}(x, Gy) & \longrightarrow & \mathcal{D}(Fx, y) \\ g \downarrow & \longrightarrow & \epsilon_y \circ Fg \end{array}$$

$$\begin{array}{ccc} Fx & \xrightarrow{Fg} & FGy \\ \downarrow & & \downarrow \epsilon_y \\ \epsilon_y \circ Fg & \rightsquigarrow & y \end{array}$$

Proof. For  $\eta_x : X \rightarrow GFx$

$$D(Fx, Fx) \cong C(X, GFx).$$
$$id_{Fx} \longrightarrow \eta_x$$

naturality of the  $\Rightarrow \eta$  is a natural transformation.

For  $\varepsilon_y : FGy \rightarrow y$

$$C(Gy, Gy) \cong D(FGy, y)$$
$$id_{Gy} \longrightarrow \varepsilon_y$$

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Ex: Products instantaneously

$$C(X, Y_1 \times Y_2) \cong C^2((X, X), (Y_1, Y_2))$$

"A map  $X \rightarrow Y_1 \times Y_2$  is the same as a pair of maps  $X \rightarrow Y_1, X \rightarrow Y_2$ ".

Product is a right adjoint to the diagonal functor

$$1_C \longrightarrow T \Delta$$

$$X \longrightarrow X \times X$$

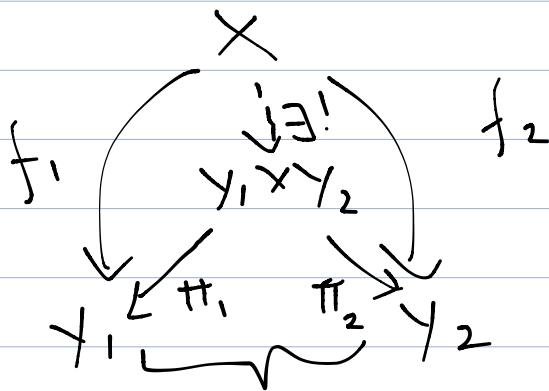
$$\begin{array}{ccc} \Delta \pi & \longrightarrow & \downarrow \rho \\ (x \times y_2, y_1 \times y_2) & \longrightarrow & (y_1, y_2) \end{array}$$

$$\eta_X: X \longrightarrow X \times X$$

$\eta_X$ : diagonal: induced by  $\text{id}: X \rightarrow X$ ,  $\text{id}: X \rightarrow X$

$\epsilon_{Y_1, Y_2}$ : projections  $\pi_1: Y_1 \times Y_2 \rightarrow Y_1$ ,

$$\pi_2: Y_1 \times Y_2 \rightarrow Y_2.$$



$$(X, X) \xrightarrow{\Sigma(f_1, f_2)} (Y_1, Y_2)$$

$$\begin{array}{ccc} & \downarrow (\Delta, \Delta) & \nearrow \exists! \\ (X \times X, X \times X) & & \end{array}$$

$$F \longrightarrow G$$

$$\eta: 1_C \longrightarrow G^F$$

$$\Sigma: FG \longrightarrow 1_D$$

$$F = F1_c \xrightarrow{1_F n} FGF \xrightarrow{\varepsilon 1_F} 1_{\mathcal{B}} F$$

$$G = 1_G G \xrightarrow{1_G \circ} GFG \xrightarrow{1_G \circ} G 1_D = G$$

id

$$\left[ \begin{array}{ccc} Fx & \xrightarrow{F(n)} & FG\bar{F}x \\ & | & | \\ & f(\cdot x) & \xrightarrow{\quad} & Fx \\ & = \text{id}_{Fx} & & \end{array} \right]$$

$\exists!$   $\sigma$   
 $F(Fx)$   
 $\epsilon_{Fx}$

$$FGFx \xrightarrow{\quad \quad \quad} Fx$$

$\downarrow F(n)$

$$\overline{FGF}x$$

Wave my hands, paste the triangles together

$$Fx \xrightarrow{F(n)} FGFx$$

$\downarrow \varepsilon_{Fx}$

$F(id_x) = id$

$$\text{Hom}(Fx, Fx) \cong \text{Hom}(x, GFx)$$

$\eta$  is the image  
of identity under  
the isomorphism

$$Fx \xrightarrow{F(n)} FGFx$$

$\downarrow \varepsilon_{Fx}$

$\varepsilon_{Fx} \circ F(n) = id_{Fx}$

$$G\gamma \xrightarrow{n_{G\gamma}} GF\gamma$$

$\downarrow G(\varepsilon_\gamma)$

$$\text{id}_{Gy} \rightarrow Gy$$

If  $F, G$  are adjoints, then it's not necessary for  $F, G$  to be inverses (equivalence-inverses). But they do satisfy composing weakly to the identity in their images.

$$G(\varepsilon_y) \circ \eta_{Gy} = \text{id}_{Gy}$$

$$\varepsilon_{Fx} \circ F(\eta_x) = \text{id}_{Fx}$$

instantiating hom-set  $\mathcal{I}_0$

Theorem: If  $F \dashv G$ ,  $J: I \rightarrow \mathcal{C}$ , and  $\mathcal{C}$  has a colimit  $\text{colim } J$ , then

$$F(\text{colim } J) = \text{colim } (F \circ J)$$

$$\zeta: J \rightarrow \Delta \text{colim } J$$

$\zeta_x: J_x \rightarrow \text{colim } J$  s.t.  
for any  $f \in I(x, x')$ ,

$$J_x \xrightarrow{Jf} J_{x'}$$

$$\begin{array}{ccc} \zeta_x & \searrow & \zeta_{x'} \\ & \text{colim } J & \end{array}$$

In  $\mathcal{D}$ , take the object  $F(\text{colim } J)$   
and the morphisms

$$\begin{array}{ccc} FJ_x & \longrightarrow & F(\text{colim } J) \\ F(\zeta_x) & & \end{array}$$

Say we have  $y \in \mathcal{D}$ ,  $\phi: FJ \rightarrow Gy$

$$\phi_x: FJ_x \rightarrow y$$

$$G\phi_x: GFJ_x \rightarrow Gy$$

$$\begin{array}{ccccc} J_x & \xrightarrow{\eta_{J_x}} & GFJ_x & \xrightarrow{G\phi_x} & Gy \\ \zeta_x \downarrow & & \downarrow GF\zeta_x & & \\ \text{colim } J & \xrightarrow{\eta_{\text{colim } J}} & GF\text{colim } J & & \end{array}$$

$G\phi_x \circ \eta_{J_x}$  is a corone  $J \rightarrow Gy$ .

By the property of  $\text{colim } J$ ,  $J'$

$\text{colim} (G \dashv \times \circ \eta_{Jx}) : \text{colim } J \rightarrow G$   
making the diagram commute.

$$\begin{array}{ccc}
 \text{colim } J & \longrightarrow & G \\
 \downarrow \eta_{\text{colim } J} & \nearrow J' \cdot g \text{ sat} & \uparrow F \text{ colim } J \rightarrow y \\
 GF \text{ colim } J & & Gg \text{ makes} \\
 & & \text{the diagram} \\
 & & \text{commute}
 \end{array}$$

$F \text{ colim } J$ ,  $F \zeta_x : FJx \rightarrow F \text{ colim } J$   
are the colimit.

So left adjoints preserve colimits, similarly,  
right adjoints preserve limits.

$F : \underline{\text{Set}} \rightarrow \underline{k\text{Vect}}$  free functor

$\sqcap$  products in Set are disjoint union  
coproducts in kVect are direct sum

$$F(\coprod_{i \in I} S_i) \cong \bigoplus_{i \in I} F(S_i)$$

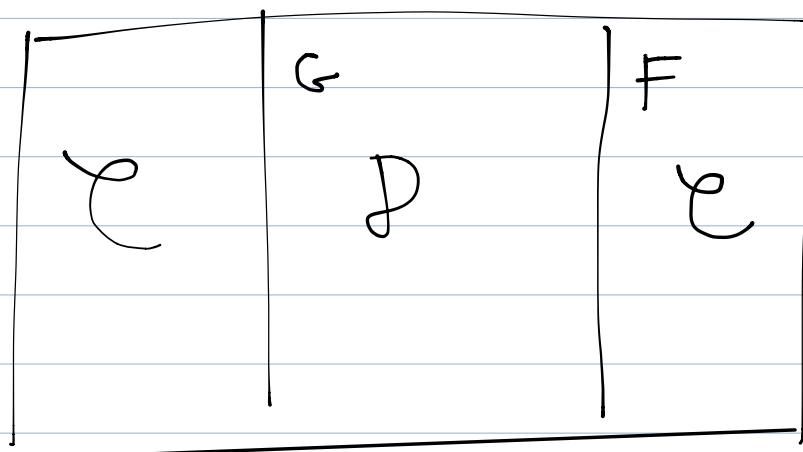
Ex.  $\mathcal{C} = \underline{\text{Top}}$ ,  $G: \text{Top} \rightarrow \text{Set}$  forgetful.

$$H \dashv G \dashv F$$

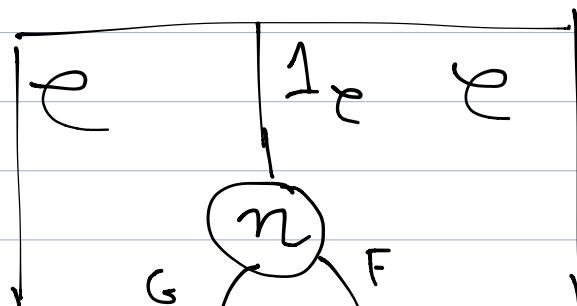
$\uparrow$  "discrete"       $\curvearrowright$  "indiscrete"

$T_G$  compute a limit/colimit in  $\underline{\text{Top}}$ , do it to the underlying set and topologize appropriately.

String diagram:

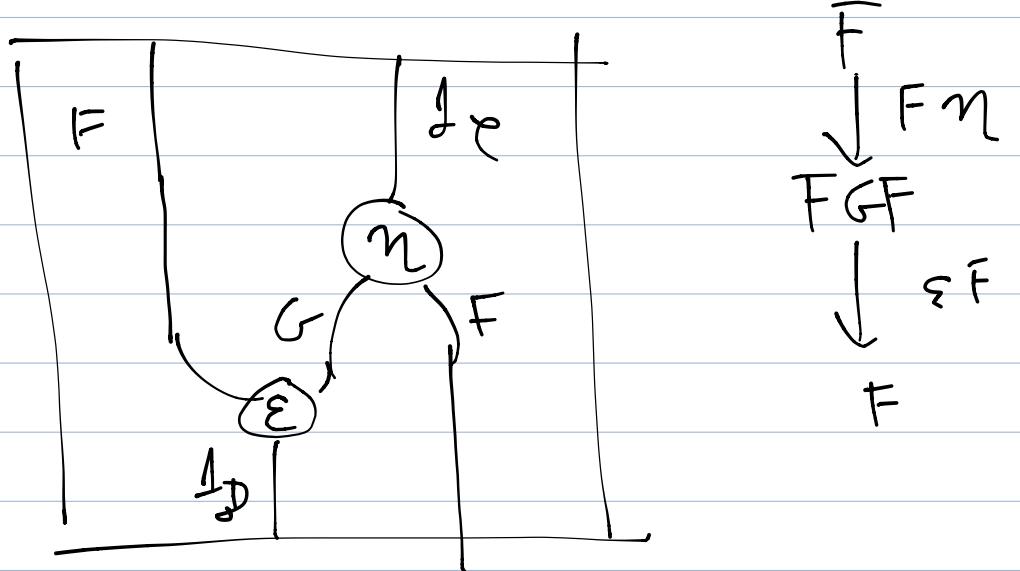


$$G \circ F : \mathcal{C} \rightarrow \mathcal{C}$$

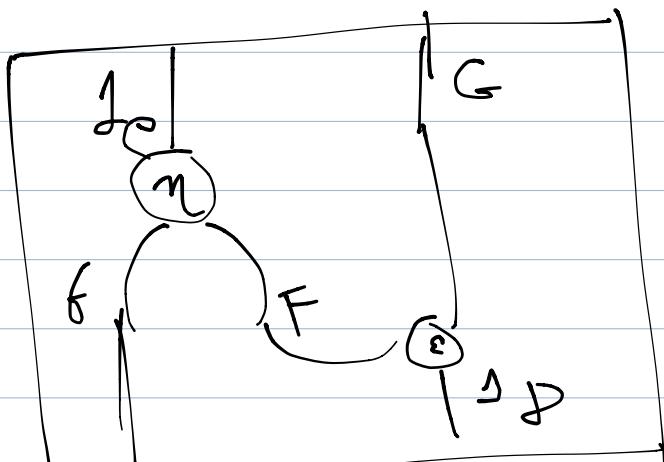
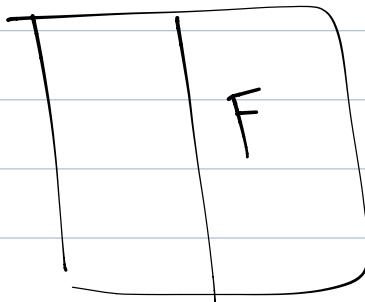


( P )

$\eta : 1_{\mathcal{C}} \longrightarrow GF$



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