

$\mathcal{C} = \text{Sets with G-actions (G-Sets)}$

$$\begin{array}{ccc} X & \xrightarrow{\rho(g)} & X \\ \phi \downarrow & & \downarrow \phi \\ Y & \xrightarrow{\sigma(g)} & Y \end{array} \quad \begin{array}{c} (X, \rho) \\ (Y, \sigma) \end{array}$$

Objects : Functors $BG \rightarrow \underline{\text{Set}}$

Morphisms : Natural transformations

Tannaka Reconstruction Theorem

G can be recovered from the category $G\text{-Set}$
and the forgetful functor $U: G\text{-Set} \rightarrow \underline{\text{Set}}$.

$$G \cong \text{Nat}(U, U)$$

BG is the category with object $*$

and $BG(*, *)$ are exactly G .

$$F: BG \rightarrow \underline{\text{Set}}$$

$$* \longmapsto X$$

$$g \longmapsto \rho(g): X \rightarrow X$$

Claim: $U \cong \text{Hom}((G \text{ left action}), -)$.

Natural $\eta: U \rightarrow \underset{G\text{-Set}}{\operatorname{Hom}}((G, H), -)$.

$\eta_{(X, f)}: U((X, f)) \rightarrow \underset{G\text{-Set}}{\operatorname{Hom}}((G, H), (X, f))$

$$X \ni x \mapsto (1 \mapsto x)$$

elements of $X \hookrightarrow G\text{-Set maps}$
 $(G, H) \rightarrow X$

1. A morphism $\phi: (G, H) \rightarrow (X, f)$ is

uniquely determined by where it sends 1:

$$g = g \cdot 1 \xrightarrow{\phi} f(g) \cdot \phi(1)$$

2. Given any $x \in X$, there is such a map

$$\delta_x: (G, H) \rightarrow (X, f).$$

$$g \mapsto f(g) \cdot x$$

Have to check naturality: $(X, f), (Y, \sigma)$

$$\phi: (X, f) \rightarrow (Y, \sigma)$$

$$x \in U(X, f) \xrightarrow{U(\phi)} U(Y, \sigma)$$

$$\eta$$

$$\eta$$

$$G\text{-Set}((G, H), (X, f)) \xrightarrow{\phi \circ -} G\text{-Set}((G, H), (Y, \sigma))$$

Right, then down:

$$x \longmapsto b(x) \longmapsto \text{the unique map } 1 \rightarrow \psi(x)$$

Down, then right:

$$x \longmapsto \text{the unique map } 1 \rightarrow x \longmapsto 1 \rightarrow \delta(x)$$

$$U \cong \underline{G\text{-Set}}((G, \text{id}), -).$$

$$\text{Nat}(U, U) \cong G\text{-Set}((G, \text{id}), (G, \text{id}))$$

Setwise, this is just $U((G, \text{id})) = G$.

$$\begin{aligned} G &\longrightarrow \text{Nat}(U, U) \\ g &\longmapsto (\eta_{(x, g)} = g) \end{aligned}$$

$$G \cong \text{Nat}(U, U).$$

$$G \not\cong \underbrace{\text{Nat}(\text{id}, \text{id})}_{\cong \mathbb{Z}(G)}$$

elements of this should be G -maps

$$\begin{array}{ccc} (X, \rho) & \xrightarrow{\phi} & (Y, \sigma) \\ \circlearrowleft \eta_{(X, \rho)} \downarrow & & \downarrow \eta_{(Y, \sigma)} \circlearrowright \\ (X, \rho) & \xrightarrow{\phi} & (Y, \sigma) \end{array}$$

Should be
G-maps

Tannaka for Rings

$R\text{-Mod}$ is the additive functor category

$$\text{Ab}^R = BR \longrightarrow \text{Ab}$$

Ab -enriched functors $BR \rightarrow \text{Ab}$
 $F(r+s) = F(r) \sqcup F(s)$

$$R \cong M \text{ then } (r+s) \cdot m = r \cdot m + s \cdot m$$

Let $U: R\text{-Mod} \rightarrow \text{Ab}$ be the forgetful functor. Then $R \cong \text{Nat}(U, U)$.

$$r \mapsto (\eta_M(x) = r \cdot x).$$

$$U \cong \text{Hom}_R(R, -).$$

(A map $R \rightarrow M$ is uniquely determined by where it sends 1 and for every element,

there is a map $\mathbb{Z} \rightarrow M^1$.

$$\text{Hom}_R(RR, RR) \cong R.$$

$$\begin{array}{ccccc}
 & & x \mapsto sx & & \\
 \downarrow & R & \xrightarrow{\quad} & R & \downarrow \\
 \checkmark & \downarrow & & \downarrow & \text{right multi} \\
 \checkmark & R & \longrightarrow & R & \checkmark \quad \text{by } r \\
 & & x \mapsto sx & &
 \end{array}$$

$$Sr = Sr.$$

$$U = \text{Hom}_R(RR, -)$$

$$\text{Nat}(U, U) \cong \text{Hom}_R(RR, RR).$$

Non-Tannaka

$\mathcal{C} = k\text{-Vect}$

$$U : \mathcal{C} \longrightarrow \text{Ab}$$

$$V \longmapsto V \otimes k^2$$

$$U = - \otimes k^2$$

U is additive faithful: $f, g : V \rightarrow W$,
 $f \otimes \text{id}_{k^2} = g \otimes \text{id}_{k^2} \iff f = g$.

$$\text{Nat}(U, U) = M_2(k).$$

The categories $k\text{-Vect}$ and $M_n(k)\text{-Mod}$
 are equivalent. For any k , there is
 an additive faithful functor $F_k : k\text{-Vect} \rightarrow$
~~is st~~

$$\text{Nat}(F_k, F_k) \cong M_n(k).$$

$$\text{Nat}(\text{id}, \text{id}) \cong \mathcal{Z}(R),$$

$$\mathcal{Z}(M_n(k)) \cong k.$$

$$\mathcal{Z}(M_m(k)) \cong \mathcal{Z}(M_m(k)).$$

$$G_1 = \mathbb{Z}_2 \times \mathbb{Z}_2, G_2 = \mathbb{Z}_4.$$

$$\text{Rep}_{\mathbb{C}}(G_1), \text{Rep}_{\mathbb{C}}(G_2)$$



$$[BG, \mathbb{C}\text{-Vect}] \quad [BG_2, \mathbb{C}\text{-Vect}]$$

" \mathbb{C} -linear representations"

$$\text{Rep}_{\mathbb{C}}(\mathbb{Z}_2 \times \mathbb{Z}_2) \cong \text{Rep}_{\mathbb{C}}(\mathbb{Z}_4)$$

S1

$$\mathbb{C}\text{-Vect} \times \mathbb{C}\text{-Vect} \times \mathbb{C}\text{-Vect} \times \mathbb{C}\text{-Vect}$$

$\mathbb{C}\text{Vect} \cong \mathbb{C}\text{Vect}^+ \cong \text{Rep}_{\mathbb{C}}(G)$

$\mathbb{C}\text{Vect}^+$

$$\begin{array}{ccc} G & \xrightarrow{i} & \mathbb{C}G \\ & \searrow & \downarrow \exists! \\ & & \text{End}_{\mathbb{C}}(V) \end{array}$$

$$\mathbb{C}G = \left\{ \sum_{\text{finite}} a_g g \mid g \in G \right\}$$

$$g \cdot h = gh, \text{ extend linearly}$$

$$\text{Rep}_{\mathbb{C}}(G) \cong \mathbb{C}G\text{-Mod.}$$

$$\mathbb{C}(\mathbb{Z}_2 \times \mathbb{Z}_2), \quad \mathbb{C}(\mathbb{Z}_4) \cong \mathbb{C}^4.$$

What if there's no rep. object for the forgetful?

\mathcal{C} is unitary reps of $G = \mathbb{Z}$:

$1 \mapsto$ an operator on a Hilbert space
 H with $T^*T = TT^* = 1$.

$\text{Nat } (\mathbb{N}, \mathbb{N}) \cong \widehat{\mathbb{Z}},$ profinite completion.

