## Finite Element Methods: Numerical Exercises Paul J. Atzberger

- 1. Show that each of the elements have the stated regularity as follows:
  - (a) Lagrange triangular element based on  $\mathcal{P}_k$  with k + 1 distinct nodes along each edge is  $C^0$ .
  - (b) Hermite triangular element based on  $\mathcal{P}_3$  is  $C^0$ .
  - (c) Argyris triangular element based on  $\mathcal{P}_5$  is  $C^1$  in the normal direction across edges.



Figure 1: Triangular Elements.

2. Consider the elliptic PDE (Poisson problem) given by

$$\Delta u(\mathbf{x}) = -f(\mathbf{x}), \ \mathbf{x} \in \Omega, \quad u(\mathbf{x}) = 0, \ \mathbf{x} \in \partial \Omega,$$

where  $\Omega = [-L, L] \times [-L, L] \subset \mathbb{R}^2$ . In the Ritz-Galerkin approximation, we seek a solution  $u_h \in \mathcal{V}_h \subset \mathcal{V} = H_0^1(\Omega)$  with

$$a(u_h, w) = -\langle f, w \rangle_{L^2}, \ \forall w \in \mathcal{V}_h,$$

where  $a(u_h, w) = \int_{\Omega} \nabla_{\mathbf{x}} u_h(\mathbf{x}) \cdot \nabla_{\mathbf{x}} w(\mathbf{x}) d\mathbf{x}$  and  $\langle f, w \rangle_{L^2} = \int_{\Omega} f(\mathbf{x}) w(\mathbf{x}) d\mathbf{x}$ . Consider a basis of functions  $\{\phi_k\}_{k=1}^N$  for  $\mathcal{V}_h$ . We can represent any  $v \in V_h$  by  $v(\mathbf{x}) = \sum_i v_i \phi_i(\mathbf{x}), u_h(\mathbf{x}) = \sum_i u_i \phi_i(\mathbf{x})$ , and approximate f by  $f_h(\mathbf{x}) = \sum_i f_i \phi_i(\mathbf{x})$ . The FEM approximation  $u_h$  can be expressed as solving the linear system

$$A\mathbf{u} = -M\mathbf{f}.$$

The A is the stiffness matrix given by  $A_{ij} = a(\phi_i, \phi_j)$ , M is the mass matrix given by  $M_{ij} = \langle \phi_i, \phi_j \rangle_{L^2}$ , and  $[\mathbf{u}]_i = u_i$ ,  $[\mathbf{f}]_i = f_i$ .

(a) (Meshing) Discretize the domain  $\Omega$  into elements  $\mathcal{T} = {\mathcal{T}_{\ell}}_{\ell=1}^{m}$ , where  $\mathcal{T}_{\ell}$  are triangular elements. For the square domain  $\Omega = [-L, L] \times [-L, L] \subset \mathbb{R}^2$ , one way to discretize is to define a coarse mesh. A basic algorithm to obtain a more refined discretization is to loop over each triangle and bisect the edges to obtain four smaller triangles, see Figure 2. Data structures for this are a list of vertices  $\mathbf{v}_i \in \mathbb{R}^2$  and tuples  $(i_1, i_2, i_3)$  which give the indices of the vertices of each triangle.



Figure 2: Mesh triangulation and refinement by triangle bisection.

Implement this meshing algorithm for the triangulation in Figure 2. Plot the triangulations when this refinement procedure is done up to n = 5 times.

(b) (Assembly and Quadratures) For the discretization into triangular elements  $\mathcal{T} = \{\mathcal{T}_{\ell}\}_{\ell=1}^{m}$ , take  $\{\phi_k\}_{k=1}^{N}$  to be the nodal basis functions for Lagrange elements with polynomial shape functions of degree dso that  $v_h | \mathcal{T}_{\ell} \in \mathcal{P}_d$ . The stiffness matrix A is obtained through an assembly procedure where we compute the integral by breaking it into parts summing up the inner-products over each element  $\mathcal{T}_{\ell}$  as  $A_{ij} = a(\phi_i, \phi_j) = \sum_{\ell=1}^{m} \int_{\mathcal{T}_{\ell}} \nabla_{\mathbf{x}} \phi_i(\mathbf{x}) \cdot \nabla_{\mathbf{x}} \phi_j(\mathbf{x}) d\mathbf{x} = \sum_{\ell=1}^{m} A_{\ell,ij}$ , and similarly,  $M_{ij} = \langle \phi_i, \phi_j \rangle_{L^2} = \sum_{\ell=1}^{m} \int_{\mathcal{T}_{\ell}} \phi_i(\mathbf{x}) \phi_j(\mathbf{x}) d\mathbf{x} = \sum_{\ell=1}^{m} M_{\ell,ij}$ . Integrals are approximated by high-precision quadratures

$$\tilde{A}_{\ell,ij} = \sum_{k} \omega_k \nabla_{\mathbf{x}} \phi_i(\mathbf{x}_k) \cdot \nabla_{\mathbf{x}} \phi_j(\mathbf{x}_k), \quad \tilde{M}_{\ell,ij} = \sum_{k} \omega_k \phi_i(\mathbf{x}_k) \phi_j(\mathbf{x}_k).$$

The  $\{\omega_k\}$  are the quadrature weights and  $\{\mathbf{x}_k\}$  are the quadrature nodes. Note in general the quadrature nodes can differ from the

finite element nodes. We use these approximations to obtain

$$\tilde{A}\mathbf{u} = -\tilde{M}\mathbf{f}.$$

For the case of Lagrange elements using polynomial spaces of degree d, we use quadratures that have order 2d. This allows for computing the integrals up to round-off errors. For quadratures on triangulations, see Figure 3 and Table 1.



Figure 3: Quadrature Nodes.

d	n	k	$\mathbf{x}_k$	$\omega_k$	k	$\mathbf{x}_k$	$w_k$	k	$\mathbf{x}_k$	$w_k$	k	$\mathbf{x}_k$	$w_k$
1	1	1	(1/3, 1/3)	1/2									
2	3	1	(1/6, 1/6)	1/6	2	(2/3,1/6)	1/6	3	(1/6, 2/3)	1/6			
3	4	1	(1/3, 1/3)	-9/32	2	(3/5,1/5)	25/96	3	(1/5, 3/5)	25/96	4	(1/5, 1/5)	25/96
4	7	1	(0,0)	1/40	2	(1/2,0)	1/15	3	(1,0)	1/40			
		4	(1/2,1/2)	1/15	5	(0,1)	1/40	6	(0,1/2)	1/15	7	(1/3, 1/3)	9/40

Table 1: Quadratures on triangulations for  $\int_0^1 \int_0^{1-x_1} f(\mathbf{x}) d\mathbf{x} \approx \sum_k f(\mathbf{x}_k) w_k$ ,  $\mathbf{x} = (x_1, x_2)$ . The *d* is the quadrature order, *n* number of nodes,  $\mathbf{x}_k$  nodes, and  $\omega_k$  weights. For affine reference element map  $\mathbf{x} = \psi(\mathbf{X})$  with  $\psi(\mathcal{T}_\ell) = \mathcal{T}_0$  and Jacobian  $J(\mathbf{X}) = |\det \partial \psi / \partial \mathbf{X}|$ , the quadrature is applied using  $\int_{\mathcal{T}_\ell} F(\mathbf{X}) d\mathbf{X} = \int_{\mathcal{T}_0} F(\psi^{-1}(\mathbf{x})) J^{-1} d\mathbf{x}$ .

Using this assembly + quadrature approach, implement codes to compute for a given triangulation the stiffness and mass matrices when d = 1 and d = 2.

Consider the FEM approximation for the solutions u with  $L = \pi$  and (i)  $u(x_1, x_2) = \cos(5x_1) \sin(5x_2)$  and (ii)  $u(x_1, x_2) = \exp(-\cos(3x_1) + \sin(3x_2))$ . Use  $f(\mathbf{x}) = -\Delta u$  evaluated at the nodal points to obtain the numerical data for these test problems.

Make a log-log plot of the solution error vs mesh size  $h^{-1} = 2^{-n}$  for meshes with refinements n = 1, 2, ..., 5. What is the exhibited order of accuracy of the Lagrange FEMs when d = 1 and d = 2? (c) (Iterative Methods) To solve approximately

 $A\mathbf{u} = \mathbf{b}$ , where  $\mathbf{b} = -M\mathbf{f}$ ,

iterative methods can be used of the form

$$B\mathbf{u}^{n+1} = C\mathbf{u}^n + \mathbf{b}$$

For convergence, B-C = A and the spectral radius of  $B^{-1}C$  is taken to satisfy  $\rho(B^{-1}C) < 1$ . It is common to decompose the matrix as A = D - L - U, where D is the diagonal entries, -L the lower entries, and -U the upper entries. A few example iterative methods are

- i. Direct Relaxation with B = I and  $C = I + \eta A$ , with small enough  $\eta$  s.t.  $\eta \leq 2/\lambda$  or smaller, where  $\lambda$  is the largest eigenvalue of A.
- ii. Jacobi Iteration with B = D and C = L + U.
- iii. Gauss-Seidel Iteration with B = D + L and C = U.

Compare these methods for approximating the solution **u** when  $L = \pi$ and (i)  $u(x_1, x_2) = \cos(5x_1)\sin(5x_2)$  and

(ii)  $u(x_1, x_2) = \exp(-\cos(3x_1) + \sin(3x_2))$ . Use  $f(\mathbf{x}) = -\Delta u$  evaluated at the nodal points to obtain the numerical data for these test problems.

Make a log-log plot of the number iterations and the error for meshes with n = 5 refinements. How many iterations does each method need to converge to 1% accuracy for solving the linear system? We remark that in practice these convergence rates are further enhanced by using preconditioners.