

Homework 2 Solutions

CS117

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① Define the propositions $\{Q_1, Q_2, \dots\} = \{Q_\ell : \ell \in \mathbb{N}\}$ by

$$Q_\ell := P_{\ell-1+m}$$

Then Q_1 is true and for all $\ell \in \mathbb{N}$, $\ell-1+m \geq m$, so if $Q_\ell = P_{\ell-1+m}$ is true, then $Q_{\ell+1} = P_{(\ell-1+m)+1}$ is true. Therefore, by induction $\{Q_1, Q_2, \dots\} = \{P_m, P_{m+1}, \dots\}$ are all true.

(2) (a) Base case: When $n=2$, we have $2^2=4 > 3=2+1$.
Inductive step: Suppose $n^2 > n+1$. Then
 $(n+1)^2 = n^2 + 2n + 1 > (n+1) + 2n + 1 \geq (n+1) + 1$,
which completes the proof.

(b) Base case: When $n=4$, we have $4! = 24 > 16 = 4^2$.
Inductive step: Suppose $n! > n^2$. Then
 $(n+1)! = (n+1) \cdot n! > (n+1)n^2$. By part (a),
 $n^2 > n+1$ for all $n \geq 2$. Therefore, we
may continue the previous inequality to
obtain $(n+1)n^2 > (n+1)^2$. Hence $(n+1)! > (n+1)^2$,
which completes the proof.

③ Assume for the sake of contradiction that $q > p$. By the proposition from class, there exists $r \in F$ so that $q > r > p$.
By assumption, since $r > p$, we must have $q \leq r$. This contradicts the fact that $q > r$. Therefore, we must have $q \leq p$.

⑤

① Let $m = \max(S)$. By definition of maximum, $s \leq m$ for all $s \in S$, so m is an upper bound of S . If \tilde{m} is an upper bound of S , then since the definition of maximum ensures $m \in S$, we have $\tilde{m} \geq m$. Thus, m is the least upper bound of S , so $m = \sup(S)$.

② By definition, $\sup(S) \geq s$ for all $s \in S$. Thus, if $\sup(S) \in S$, it is the largest element of S and $\max(S) = \sup(S)$.

③ Suppose S has a maximum M_0 . By ①, $M_0 = \sup(S)$. This contradicts the fact that $\sup(S) \notin S$. Therefore, S must not have a maximum.

⑥a Let s be an element of S . Since $\inf(S)$ is a lower bound for S , $\inf(S) \leq s$. Since $\sup(S)$ is an upper bound for S , $s \leq \sup(S)$. Therefore, $\inf(S) \leq \sup(S)$.

⑥b We will show $S = \{\inf(S)\}$, so there is one element in the set. Since $\inf(S) = \sup(S)$, $\inf(S)$ is both an upper and lower bound for S . In particular, for any $s \in S$, $\inf(S) \leq s$ and $\inf(S) \geq s$. Thus $\inf(S) = s$ for all $s \in S$. This shows $S = \{\inf(S)\}$.

⑦ (a) Since $\inf(T)$ is a lower bound for the set T , if $S \subseteq T$, then $\inf(T)$ is also a lower bound for the set S . Since $\inf(S)$ is the greatest lower bound of S , $\inf(T) \leq \inf(S)$. The fact that $\sup(S) \leq \sup(T)$ follows from an analogous argument. The fact that $\inf(S) \leq \sup(S)$ follows from (6)

you should write this argument or use $-S$ and $-T$ to reduce to what you have already shown

(b) Since $\sup(S)$ is an upper bound for S and $\sup(T)$ is an upper bound for T , $\max\{\sup(S), \sup(T)\}$ is an upper bound for $S \cup T$. Thus, since $\sup(S \cup T)$ is the least upper bound for $S \cup T$, $\sup(S \cup T) \leq \max\{\sup(S), \sup(T)\}$.

By Q7 (a), since $S \subseteq S \cup T$ and $T \subseteq S \cup T$, we have $\sup(S) \leq \sup(S \cup T)$ and $\sup(T) \leq \sup(S \cup T)$. Thus $\max\{\sup(S), \sup(T)\} \leq \sup(S \cup T)$.

Combining these two inequalities, we conclude $\max\{\sup(S), \sup(T)\} = \sup(S \cup T)$.

⑧ (a) Since S is bounded below, $\exists m_0 \in \mathbb{R}$ s.t. $s \geq m_0 \forall s \in S$. This implies $-m_0 \geq -s \forall s \in S$, so $-S$ is bounded above.

(b) Since S is nonempty, so is $-S$. Since $-S$ is bounded above, by definition of the real numbers, it has a supremum, $\sup(-S)$.

(c) Since $\sup(-S)$ is an upper bound for $-S$, $-s \leq \sup(-S)$ for all $s \in S$, hence $s \geq -\sup(-S)$ for all $s \in S$. Therefore $-\sup(-S)$ is a lower bound for S , and it suffices to show it is the greatest lower bound.

for the sake of contradiction that
Suppose m_0 is a lower bound for S with $m_0 > -\sup(-S)$. As argued in part (a), $-m_0$ is an upper bound for $-S$. Furthermore, $m_0 > -\sup(-S)$ implies $-m_0 < \sup(-S)$. This is a contradiction, since $\sup(-S)$ is the least upper bound for $-S$.

Therefore $-\sup(-S)$ must be the greatest lower bound of $-S$.

⑨ Step 1: Show that for all $t \in T$, $\inf(S+T) - t$ is a lower bound for S .

By defn of $S+T$ and the infimum, $\inf(S+T)$ is a lower bound for $S+T$, so $s+t \geq \inf(S+T) \Leftrightarrow s \geq \inf(S+T) - t$ for all $s \in S, t \in T$. Thus, for all $t \in T$, $\inf(S+T) - t$ is a lower bound for S .

Step 2: Show that $\inf(S+T) - \inf(S)$ is a lower bound for T .

By Step 1, for all $t \in T$, $\inf(S+T) - t$ is a lower bound for S . By defn, $\inf(S)$ is the greatest lower bound of S .
Thus, $\inf(S) \geq \inf(S+T) - t$.
 $\Leftrightarrow t \geq \inf(S+T) - \inf(S)$ for all $t \in T$.

Since $\inf(S+T) - \inf(S)$ is a lower bound for T and $\inf(T)$ is the greatest lower bound,

$$\inf(T) \geq \inf(S+T) - \inf(S).$$

$$\inf(S) + \inf(T) \geq \inf(S+T). \quad (*)$$

It remains to prove the opposite inequality. Since $\inf(S)$ and $\inf(T)$ are lower bounds for S and T , for all $s \in S$ and $t \in T$, $\inf(S) \leq s$ and $\inf(T) \leq t \Rightarrow \inf(S) + \inf(T) \leq s + t$. Thus, $\inf(S) + \inf(T)$ is a lower bound for $S+T$. Since $\inf(S+T)$ is the greatest lower bound,

$$\inf(S) + \inf(T) \leq \inf(S+T). \quad (**)$$

Thus, combining inequalities $(*)$ and $(**)$, we obtain

$$\inf(S) + \inf(T) = \inf(S+T). \quad \square$$

- ⑩ Throughout, we use S to denote the set under consideration.
- Ⓐ $\sup(S) = \sqrt{2}$, $\inf(S) = -\sqrt{2}$
 - Ⓑ $\sup(S) = \pi$, $\inf(S) = -1$
 - Ⓒ $\sup(S) = \inf(S) = 1$
 - Ⓓ S is not bounded above, $\inf(S) = 1$
 - Ⓔ $\sup(S) = 1$, $\inf(S) = 0$

$$(f) \sup(S) = 1; \inf(S) = -1$$

$$(g) S = [-1, 1], \text{ so } \sup(S) = 1 \text{ and } \inf(S) = -1$$

(11) The flaw occurs in the final sentence. The fact that $k, m \in \mathbb{N}$ does not guarantee $k-1, m-1 \in \mathbb{N}$, as was used in the base case. For example, take $k=2, m=1$.