

MODULE CATEGORIES OVER CATEGORIES OF MODULES OVER HOPF ALGEBRAS

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Category Theory

Category theory is a branch of mathematics that studies mathematical objects by focusing on the maps between them. The most basic notion in category theory is a **category**, which consists of the following:

- A class of **objects**
- A class of **morphisms**, each of which goes between two objects
- An associative **composition** operation $\circ$ that takes a morphism $a \rightarrow b$ and a morphism $b \rightarrow c$ to obtain a morphism $a \rightarrow c$.
- Identity morphisms id_c on every object c satisfying $f \circ \text{id}_c = \text{id}_d \circ f = f$ for all morphisms $f : c \rightarrow d$.

For example, in the **category of sets**, objects are sets, and morphisms are functions between them. Composition of morphisms is just composition of functions.

An **isomorphism** is a morphism that has a left and right inverse, i.e. a morphism f such that there exists a morphism g such that $f \circ g = \text{id}$ and $g \circ f = \text{id}$. If there exists an isomorphism $f : a \rightarrow b$, we say a and b are **isomorphic**. For the purposes of category theory, isomorphic objects are the same in every way that we care about.

A **functor** $F : \mathcal{C} \rightarrow \mathcal{D}$ is a map that sends objects of $\mathcal{C}$ to objects of $\mathcal{D}$, and sends each morphism $f : a \rightarrow b$ in $\mathcal{C}$ to a morphism $F(f) : F(a) \rightarrow F(b)$ in $\mathcal{D}$. In order to be a functor, it should preserve composition, meaning $F(f \circ g) = F(f) \circ F(g)$.

If $F, G : \mathcal{C} \rightarrow \mathcal{D}$ are functors, a **natural transformation** η from F to G is a collection of morphisms $\eta_a : F(a) \rightarrow G(a)$ such that for all morphisms $f : a \rightarrow b$, the following diagram commutes:

$$\begin{array}{ccc} F(a) & \xrightarrow{\eta_a} & G(a) \\ \downarrow F(f) & & \downarrow G(f) \\ F(b) & \xrightarrow{\eta_b} & G(b). \end{array}$$

A diagram **commutes** if any two paths with matching start and end points that give you the same morphism after composing. In this case, the statement that the above diagram commutes is saying that $G(f) \circ \eta_a = \eta_b \circ F(f)$.

If each morphism η_a is an isomorphism, then we say that η is a **natural isomorphism** between F and G .

Monoidal Categories and Module Categories

To understand monoidal categories, consider the Cartesian product in the category of sets. It's almost associative, but not quite: technically $A \times (B \times C)$ contains tuples of the form $(a, (b, c))$ whereas $(A \times B) \times C$ contains tuples of the form $((a, b), c)$. While these aren't equal, they are naturally isomorphic as sets, i.e. $\times$ is associative up to natural isomorphism. This idea leads to the following definition:

A **monoidal category** is a category $\mathcal{C}$ with a functor $\otimes : \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$, an object $1_{\mathcal{C}} \in \mathcal{C}$, an **associator** natural isomorphism $\alpha_{a,b,c} : a \otimes (b \otimes c) \xrightarrow{\cong} (a \otimes b) \otimes c$, and **unitor** natural isomorphisms $\lambda_a : 1_{\mathcal{C}} \otimes a \xrightarrow{\cong} a$ and $\rho_a : a \otimes 1_{\mathcal{C}} \xrightarrow{\cong} a$ such that the following pentagon and triangle diagrams commute for all $a, b, c, d \in \mathcal{C}$:

$$\begin{array}{ccc} a \otimes (b \otimes (c \otimes d)) & \xrightarrow{\alpha_{a,b,c \otimes d}} & (a \otimes b) \otimes (c \otimes d) \\ \downarrow \text{id}_a \otimes \alpha_{b,c,d} & & \downarrow \alpha_{a \otimes b, c, d} \\ a \otimes ((b \otimes c) \otimes d) & \xrightarrow{\alpha_{a,b \otimes c, d}} & ((a \otimes b) \otimes c) \otimes d \\ & \uparrow \alpha_{a,b,c} \otimes \text{id}_d & \\ a \otimes ((b \otimes c) \otimes d) & \xrightarrow{\alpha_{a,b \otimes c, d}} & (a \otimes (b \otimes c)) \otimes d \end{array}$$

$$\begin{array}{ccc} a \otimes (1_{\mathcal{C}} \otimes b) & \xrightarrow{\text{id}_a \otimes \lambda_b} & a \otimes b \\ \searrow \alpha_{a, 1_{\mathcal{C}}, b} & & \nearrow \rho_a \otimes \text{id}_b \\ & (a \otimes 1_{\mathcal{C}}) \otimes b & \end{array}$$

A monoidal category is **strict** if $(a \otimes b) \otimes c = a \otimes (b \otimes c)$ and $1_{\mathcal{C}} \otimes a = a = a \otimes 1_{\mathcal{C}}$ and $\alpha_{a,b,c}$, λ_a , and ρ_a are all the identity.

A **linear category** is a category where for all objects a and b , the set of morphisms from a to b is a finite-dimensional complex vector space and composition is a bilinear map.

A **linear monoidal category** is a linear category with a monoidal structure such that $\otimes$ is a bilinear map.

A **module category** over a linear monoidal category $\mathcal{C}$ is a linear category $\mathcal{A}$ equipped with the following additional data:

- A bilinear functor $\otimes : \mathcal{C} \times \mathcal{A} \rightarrow \mathcal{A}$,
- A natural isomorphism $\lambda_a : 1 \otimes a \rightarrow a$,
- A natural isomorphism $\alpha_{c,d,a} : c \otimes (d \otimes a) \rightarrow (c \otimes d) \otimes a$.

Module categories must satisfy pentagon and triangle rules similar to those in the definition of a monoidal category.

Hopf Algebras

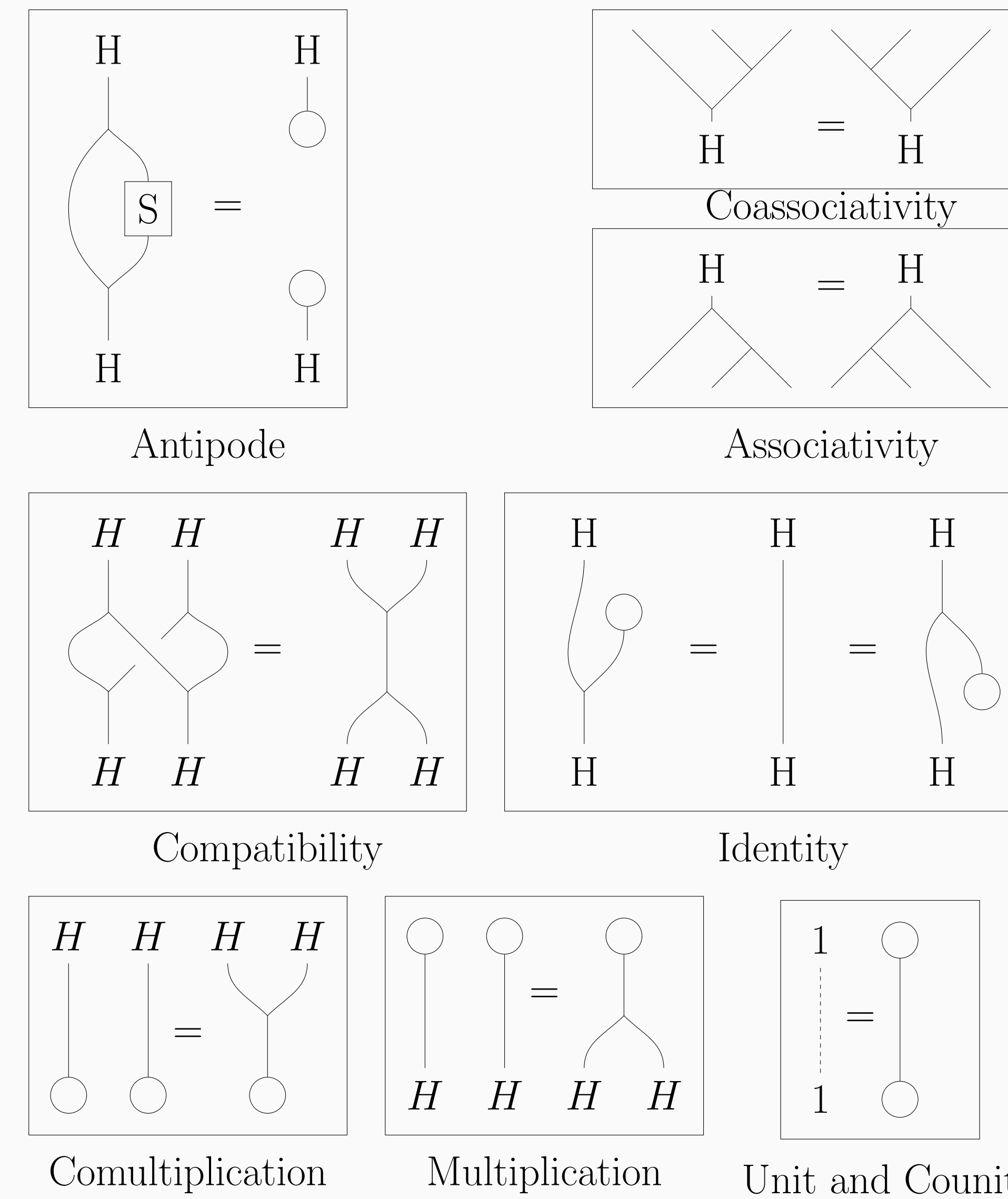
For two vector spaces V and W , their **tensor product** is the unique vector space $V \otimes W$ equipped with a bilinear map $\otimes : V \times W \rightarrow V \otimes W$ such that for any bilinear map $f : V \times W \rightarrow U$ there is a unique linear map $\tilde{f} : V \otimes W \rightarrow U$ satisfying $f(v, w) = \tilde{f}(v \otimes w)$ for all $(v, w) \in V \times W$.

An **associative algebra** over $\mathbb{C}$ is a $\mathbb{C}$ -vector space A with a multiplication operation $\cdot : A \times A \rightarrow A$ such that A forms a ring under this multiplication, and $\lambda \cdot (ab) = (\lambda a) \cdot b = a \cdot (\lambda b)$ for all $\lambda \in \mathbb{C}$ and $a, b \in A$.

A **Hopf algebra** consists of a $\mathbb{C}$ -vector space H with the following additional linear maps:

- A **multiplication** map $\nabla : H \otimes H \rightarrow H$ such that $g \cdot h := \nabla(g \otimes h)$ makes H into an associative algebra
- A **comultiplication** map $\Delta : H \rightarrow H \otimes H$
- An **antipode** map $S : H \rightarrow H$
- **Unit** and **counit** maps $\delta : \mathbb{C} \rightarrow H$ and $\varepsilon : H \rightarrow \mathbb{C}$ respectively

In order for these structures to form a Hopf algebra, we also require that they satisfy the equations expressed by the following string diagrams (and their mirror images):



Hopf Algebras continued

To interpret these string diagrams, one should read bottom to top. Each vertical strand represents the identity on H , and any two things next to each other indicate taking monoidal products. A Y-shaped branch indicates comultiplication and an upside-down Y-shape indicates multiplication. Finally, the circles represent the unit and counit respectively.

An important example of a family of Hopf algebras is the **generalized Taft algebras** $H_{n,d}(\zeta)$, parametrized by $n, d \in \mathbb{N}$ with d dividing n and $\zeta \in \mathbb{C}$ a primitive d th root of unity. As an algebra, $H_{n,d}(\zeta)$ is generated by two elements $x, g \in H_{n,d}(\zeta)$ that commute up to multiplication by ζ . As a complex vector space, $H_{n,d}(\zeta)$ is generated by the $n \cdot d$ elements $g^i h^j$ for $0 \leq i < n$ and $0 \leq j < d$, so in particular it is nd -dimensional.

Since every Hopf algebra is also an algebra and therefore a ring, we can talk about **modules** over a Hopf algebra in the usual way. Importantly, given two modules M and N over a Hopf algebra H , the tensor product of M and N as $\mathbb{C}$ -vector spaces can be given the structure of an H -module by defining $h \cdot (m \otimes n) = \sum_{i=1}^n h_i^{(1)} m \otimes h_i^{(2)} n$ where $\Delta(h) = \sum_{i=1}^k h_i^{(1)} \otimes h_i^{(2)}$. With this $\otimes$ operation, the modules over H form an example of a monoidal category $\text{Rep}(H)$.

An **algebra** in the category $\text{Rep}(H)$ consists of an H -module A with a linear map $m : A \otimes A \rightarrow A$ that is H -linear (where $A \otimes A$ has the H -module structure given above), as well as a map $\mathbb{C} \rightarrow A$, satisfying some string diagrams similar to the ones to the left. If A is an algebra in $\text{Rep}(H)$, a (right) **module** over A consists of an H -module M with a multiplication $M \otimes A \rightarrow M$ satisfying commutative diagrams similar to the ones for associativity and identity in the previous section.

The right modules over A , along with A -linear maps between them, form an example of a module category over $\text{Rep}(H)$. Furthermore, a theorem by Ostrik states that *all* module categories over $\text{Rep}(H)$ are of this form. Because of this, a large part of our research has been trying to understand what the various algebras in $\text{Rep}(H)$ are like, particularly for the case where H is a Taft algebra.

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